

November 14, 2017



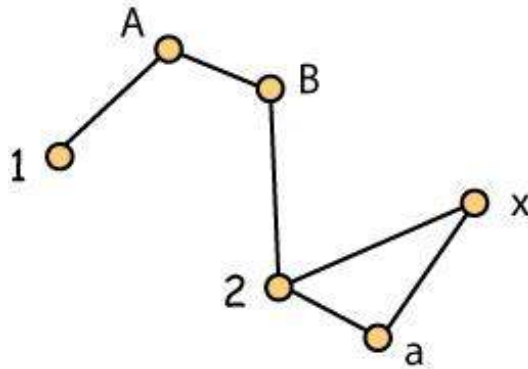
# 5 - Graph Theory Basics

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# Basic Definitions

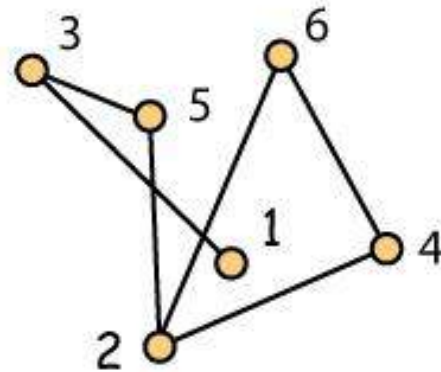
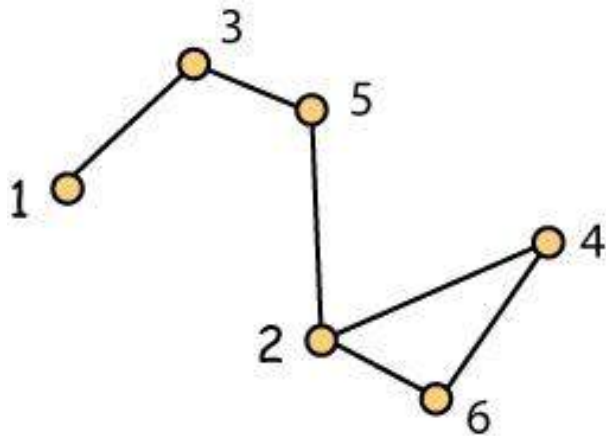
**Definition** A graph  $G$  is a pair  $(V, E)$  where  $V$  is a finite set and  $E$  is a set of 2-element subsets of  $V$ . The set  $V$  is called the **vertex** set of  $G$  and the set  $E$  is called the **edge** set of  $G$ .

**Example**  $G = (V, E)$  where  $V = \{1, 2, A, x, B, a\}$  and  $E = \{\{1, A\}, \{2, x\}, \{x, a\}, \{A, B\}, \{B, 2\}, \{2, a\}\}$ .



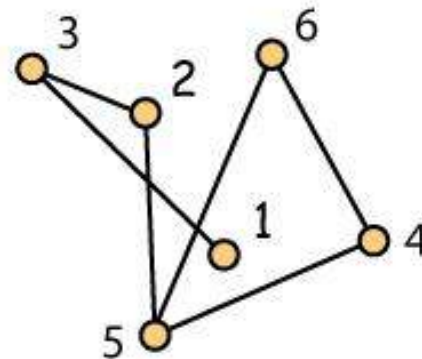
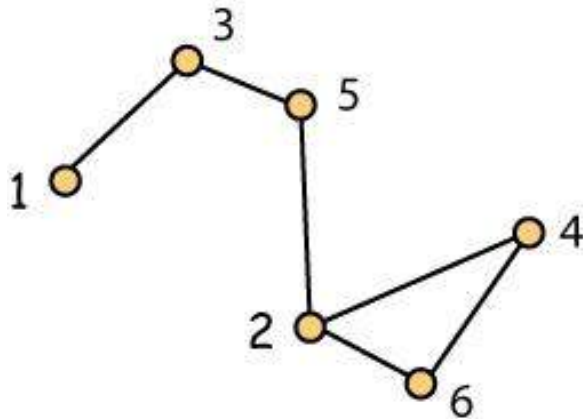
# Its All About Adjacency

**Comment** We show below two drawings of the same graph whose vertex set is  $\{1, 2, 3, 4, 5, 6\}$ .



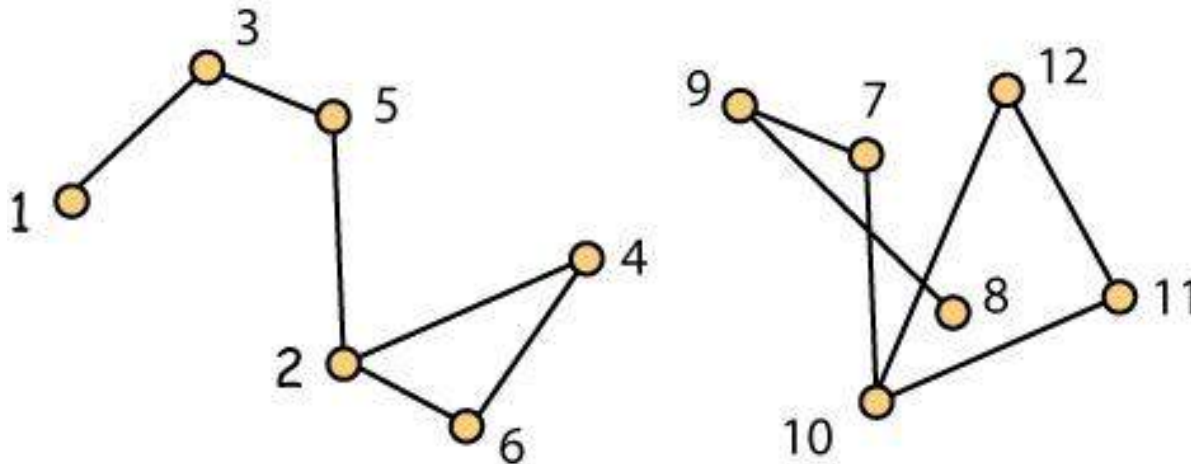
# Its All About Adjacency (2)

**Comment** We show below two drawings of graphs, each having vertex set  $\{1, 2, 3, 4, 5, 6\}$ , but now they represent different graphs.



# Its All About Adjacency (3)

**Question** Is this a drawing of one graph whose vertex set is  $\{1, 2, 3, \dots, 12\}$  or do we have drawings of two graphs, one with vertex set  $\{1, 2, 3, 4, 5, 6\}$  and the other  $\{7, 8, 9, 10, 11, 12\}$ ?



**Answer** Depends on the meaning of  $V$  in the pair  $(V, E)$ .

# Notation and Terminology

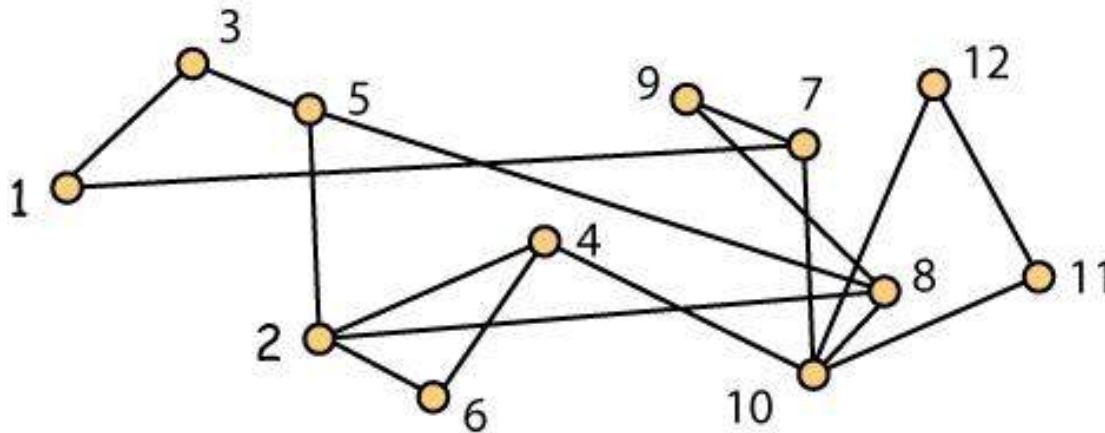
1. Vertices are also called **nodes**, **points**, **locations**, **stations**, etc.
2. Edges are also called **arcs**, **lines**, **links**, **pipes**, **connectors**, etc.
3. Remember that mathematicians are *selectively* lazy so when there is no confusion, an edge  $\{x, y\}$  will be denoted as  $xy$ . This can create some confusion when vertices are positive integers as how would one interpret a comment such as "consider the edge 2786".

# Notation and Terminology (2)

1. When  $xy$  is an edge in  $G$ , we say  $x$  and  $y$  are **adjacent** in  $G$ . Alternatively, we say they are **neighbors** in  $G$ .
2. In a graph  $G$ , the set of all neighbors of a vertex  $x$  is denoted  $N_G(x)$ . And when the graph  $G$  is fixed in the discussion, this is typically abbreviated to just  $N(x)$ .
3. The integer  $|N_G(x)|$  is called the **degree** of  $x$  in  $G$ , and is denoted  $\deg_G(x)$ . Again, when the graph is fixed, this is shortened to  $\deg(x)$ .

# Notation and Terminology (3)

**Example** A graph with vertex set  $\{1, 2, 3, \dots, 12\}$ .

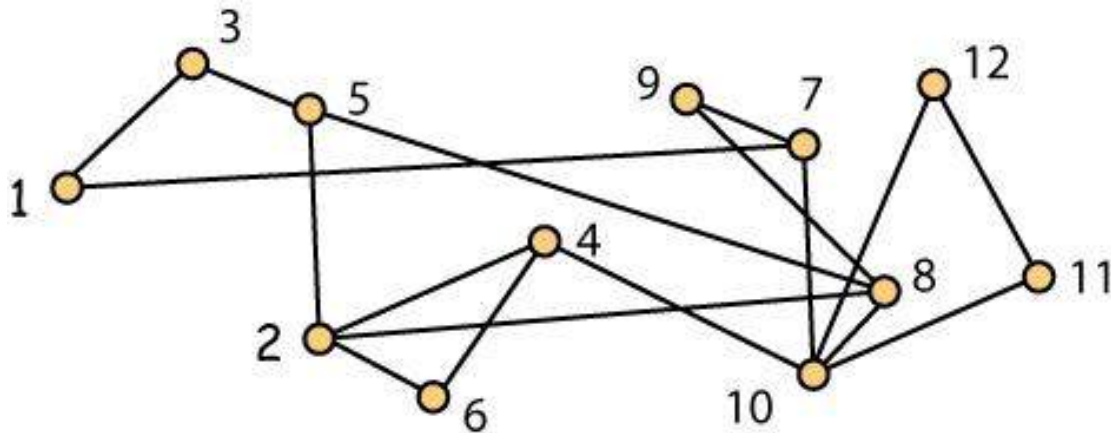


**Questions** Are 8 and 11 neighbors? What is  $\text{deg}(8)$ ?



# First Theorem in Graph Theory

**Example** Let  $G = (V, E)$  be a graph and let  $q$  be the number of edges in  $G$ . Then  $\sum_{x \in V} \deg_G(x) = 2q$



**Exercise** Verify this theorem for the graph illustrated above.

# Carlos and Dave

**Overheard in Conversation** Dave said that he was working with a graph and carefully counted all the degrees and said here is full listing of all the values:

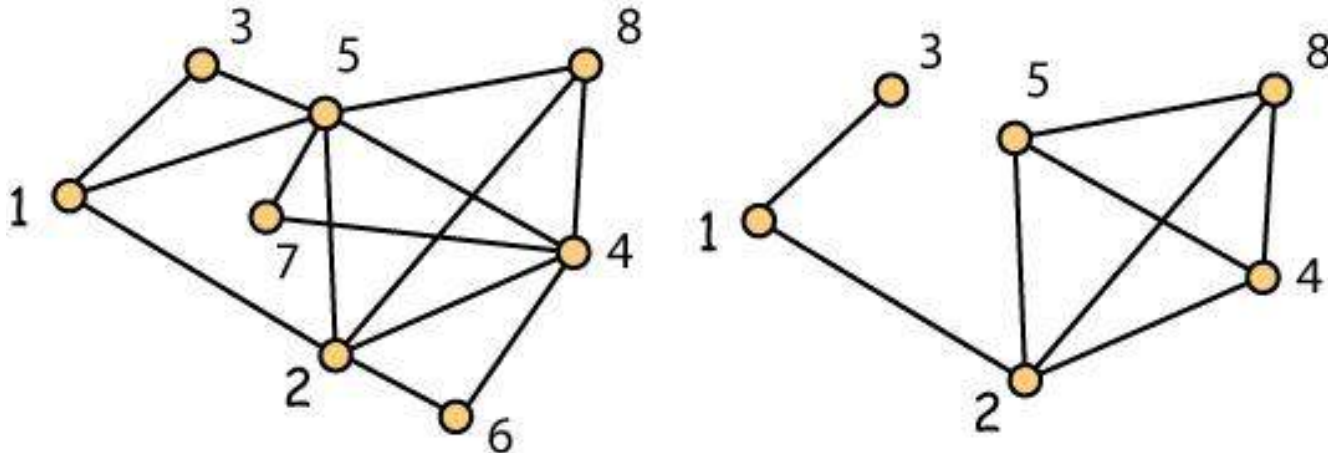
16, 13, 12, 18, 16, 22, 11, 16, 14, 10, 8, 12, 14, 16, 8, 7,  
10, 20, 12, 14, 16, 8, 6, 6, 8, 4, 8, 6, 6, 6, 6, 8, 10, 5, 8,  
8, 6, 6, 6, 6, 3, 6, 4, 8, 8, 8, 4, 8, 10, 12

Carlos remarked gently "Perhaps you should check your work."

# The Notion of a Subgraph

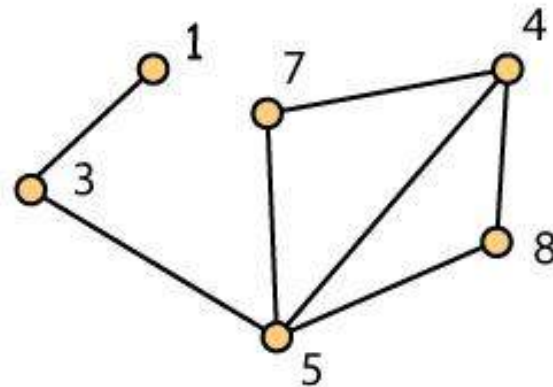
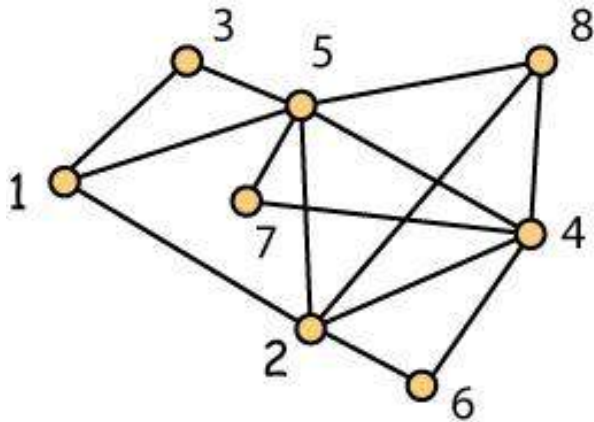
**Definition** A graph  $G' = (V', E')$  is a subgraph of a graph  $G = (V, E)$  when  $V'$  is contained in  $V$  and  $E'$  is contained in  $E$ .

**Example** On the left, we show a graph with vertex set  $\{1, 2, \dots, 8\}$ . The graph on the right is a subgraph.



# The Notion of a Subgraph

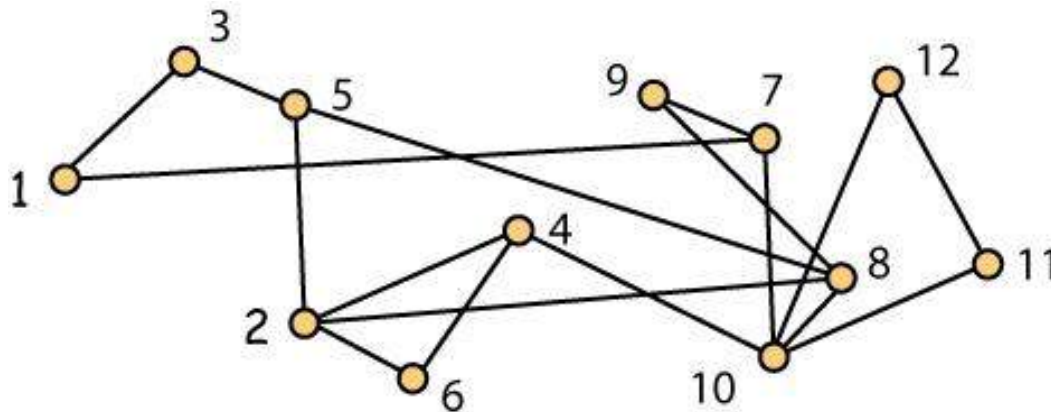
**Question** We show a graph  $G$  with vertex set  $\{1, 2, \dots, 8\}$  on the left. Is the graph on the right a subgraph?



# Paths in Graphs

**Definition** Let  $G = (V, E)$  be a graph. When  $n \geq 1$ , a sequence  $P = (x_1, x_2, \dots, x_n)$  of  $n$  distinct vertices in  $G$  is called a **path from  $x_1$  to  $x_n$  in  $G$**  if  $x_i$  is adjacent to  $x_{i+1}$  in  $G$  whenever  $1 \leq i < n$ .

**Example** In the graph shown,  $(7, 9, 8, 5, 2, 6, 4)$  is a path from 7 to 4.



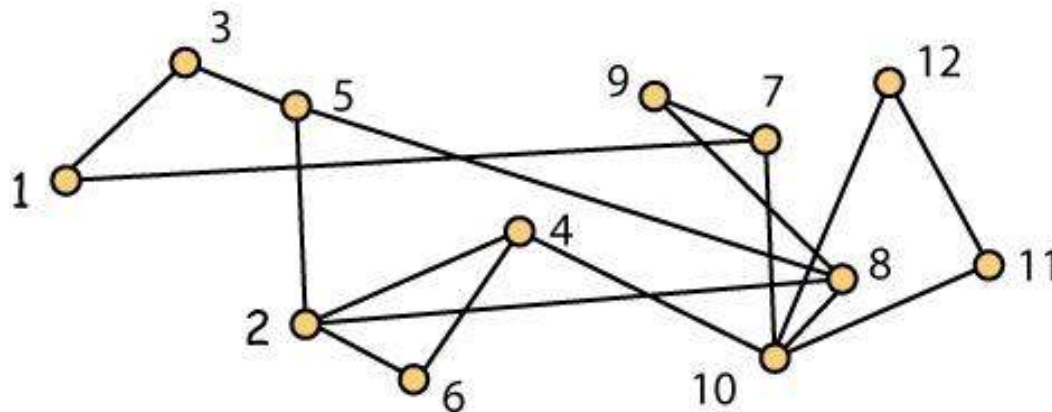
# Size of Paths

**Convention** Many authors measure how big a path is in terms of the number of edges, so they will say that a path  $(a, b, c, d, e)$  from  $a$  to  $e$  has **length** 4. In particular, they would say that when  $x$  and  $y$  are neighbors, the path  $(x, y)$  has length 1. Other authors prefer to measure paths in terms of the number of vertices, so they would say that the path  $(a, b, c, d, e)$  has **size** 5. We prefer the second option, so we will always talk about paths of a certain size and this will count the number of vertices and not the number of edges.

# Connected Graphs

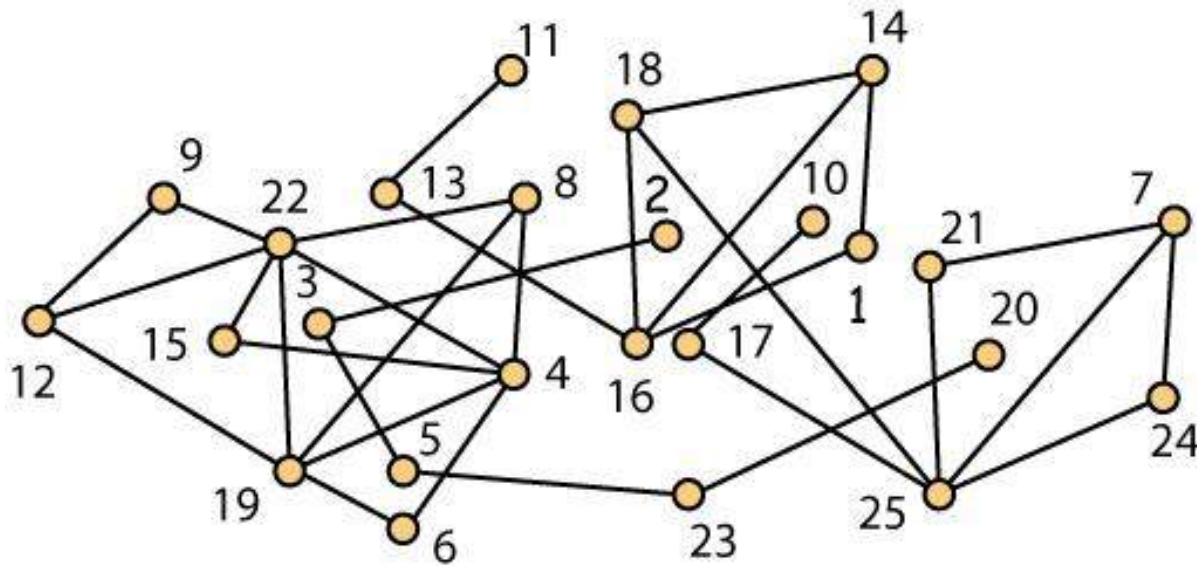
**Definition** Let  $G = (V, E)$  be a graph. We say  $G$  is **connected** if for all  $x, y$  in  $V$  with  $x \neq y$ , there is a path from  $x$  to  $y$  in  $G$ .

**Example** The graph shown below is connected.



# Connected Graphs (2)

**Definition** Let  $G = (V, E)$  be a **disconnected** graph. A subgraph  $H = (V', E')$  of  $G$  is called a **component** of  $G$  if  $H$  is connected and any subgraph of  $G$  which contains  $H$  properly is disconnected. Is this graph connected? If not, how many components does it have?

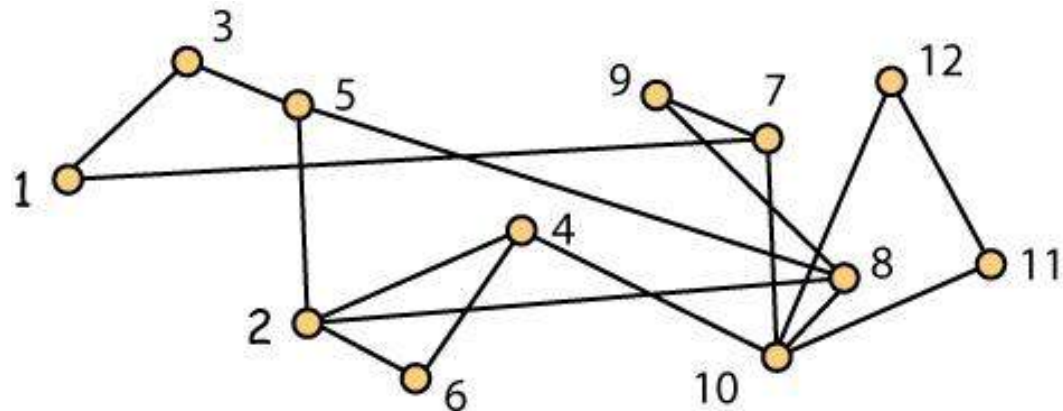




# Cycles in Graphs

**Definition** Let  $G = (V, E)$  be a graph. When  $n \geq 3$ , a sequence  $P = (x_1, x_2, \dots, x_n)$  of  $n$  distinct vertices in  $G$  is called a **cycle of size  $n$  in  $G$**  if  $x_i$  is adjacent to  $x_{i+1}$  in  $G$  whenever  $1 \leq i < n$  and  $x_n$  is adjacent to  $x_1$  in  $G$ .

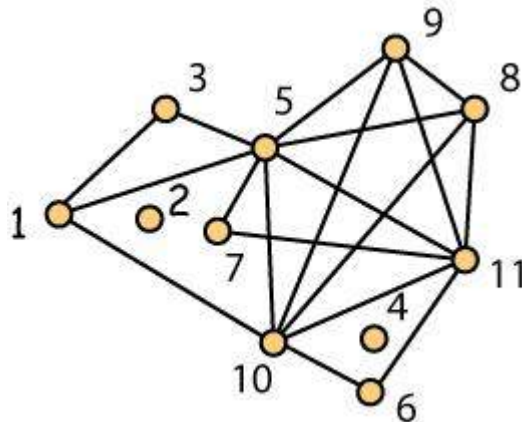
**Example** In the graph shown,  $(5, 8, 9, 7, 1, 3)$  is a cycle of size 6.



# Loose Points in Graphs

**Definition** A vertex  $x$  in a graph  $G$  is called a **loose** point (also an **isolated** point) if it has no neighbors, i.e.,  $\deg_G(x) = 0$ .

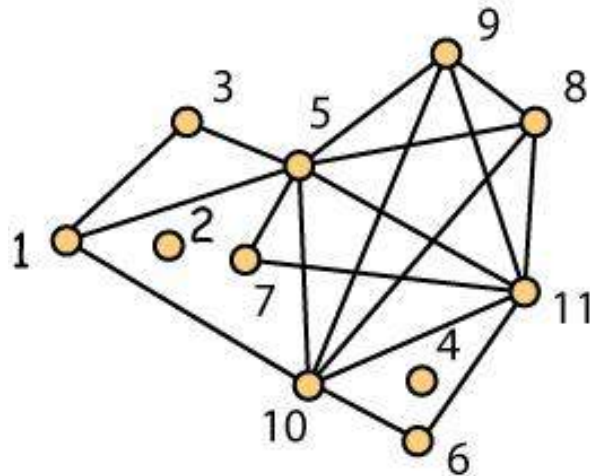
**Example** Below we show a graph with vertex set  $\{1, 2, \dots, 11\}$ . In this graph, vertices 2 and 4 are loose points.



# Cliques in Graphs

**Definition** Let  $G = (V, E)$  be a graph. When  $n \geq 1$ , a set  $S$  of vertices in  $G$  is called a **clique** if any two distinct vertices in  $S$  are adjacent in  $G$ .

**Example** In this graph, the subsets  $\{2\}$ ,  $\{6, 10\}$ ,  $\{1, 3, 5\}$  and  $\{5, 8, 9, 10, 11\}$  are cliques. There are many more.



# Xing and Zori

**Overheard in Conversations** Xing is a very good programmer and remarked to Zori that he could easily detect whether a large graph was connected and if it was disconnected whether it had any loose vertices. Zori was not impressed as she couldn't see any reason why anybody would care about either issue. Still, moderately annoyed with Xing's enthusiasm, she asked him about a problem she had read about on the web: Could he tell whether a graph on  $2n$  vertices had a clique of size  $n$ . Xing hadn't thought about it ... but now that he was challenged, he said he thought he could. Hmmmm ...

# Questions for Thought

**Challenges or Not?** Given a graph  $G = (V, E)$  with  $|V| = n$ , which of the following problems is easy and which is hard?

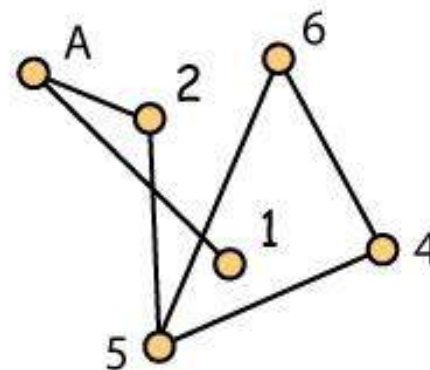
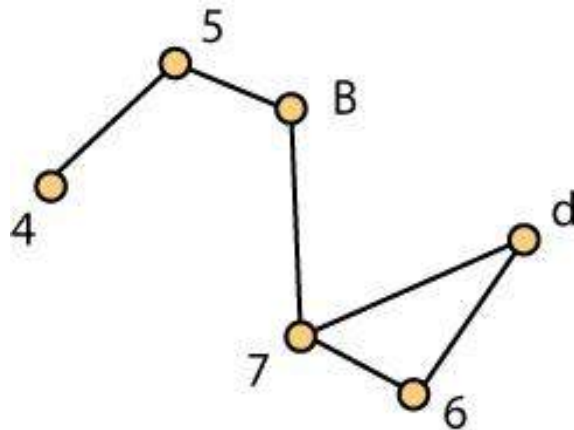
1. Is  $G$  connected?
2. Does  $G$  have a path of size at least  $n/2$ ?
3. Does  $G$  have a cycle of size at least  $n/2$ ?
4. Does  $G$  have a clique of size at least  $n/2$ ?

**Also** Suppose Alice and Bob are arguing about the correct answers to these questions for a graph with  $n = 10,000$ . Would you rather defend a "yes" answer or a "no" answer.

# Isomorphic Graphs

**Definition** Graphs  $G_1 = (V_1, E_1)$  and  $G_2 = (V_2, E_2)$  are isomorphic when there is a bijection  $f: V_1 \rightarrow V_2$  so that  $\{x, y\}$  is an edge in  $G_1$  if and only if  $\{f(x), f(y)\}$  is an edge in  $G_2$ .

**Exercise** Show that the two graphs shown below are isomorphic.

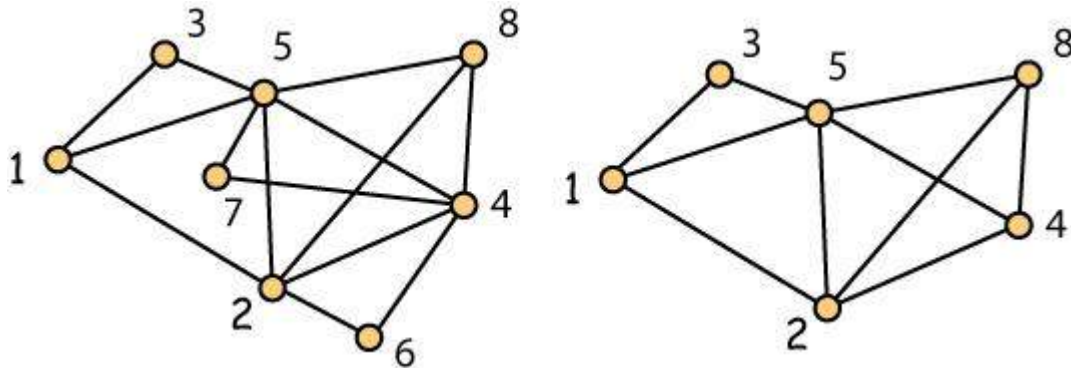


# Another Question for Thought

**Challenge or Not?** Given two graphs  $G_1 = (V_1, E_1)$  and  $G_2 = (V_2, E_2)$ , is it hard to tell whether they are isomorphic? If Yolanda says "yes" and Bob says "no", who has the easier task to convince an impartial referee?

# Induced Subgraphs

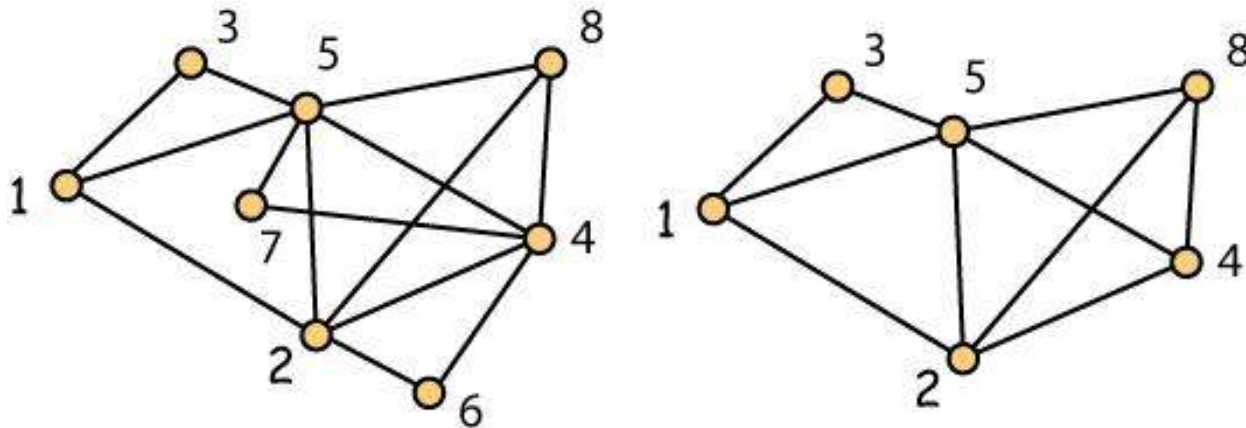
**Definition** A graph  $H = (V', E')$  is an **induced** subgraph of a graph  $G = (V, E)$  if  $V' \subseteq V$  and  $xy$  is an edge in  $H$  whenever  $x$  and  $y$  are distinct vertices in  $V'$  and  $xy$  is an edge in  $G$ . In the drawing below, the graph on the right is an induced subgraph of the graph on the left.





# Induced Subgraphs (2)

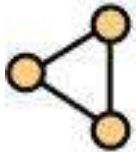
**Remark** When  $G = (V, E)$  is a graph, an induced subgraph of  $G$  is determined entirely by its vertex set, so for example, the induced subgraph on the right can be denoted as  $G - \{6, 7\}$ .



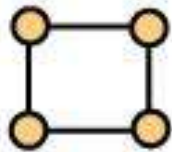


# Special Classes of Graphs

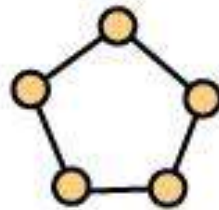
**Definition** For  $n \geq 3$ ,  $C_n$  denotes a **cycle** on  $n$  vertices. Here are  $C_3$ ,  $C_4$ ,  $C_5$  and  $C_6$ .



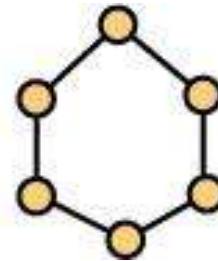
$C_3$



$C_4$



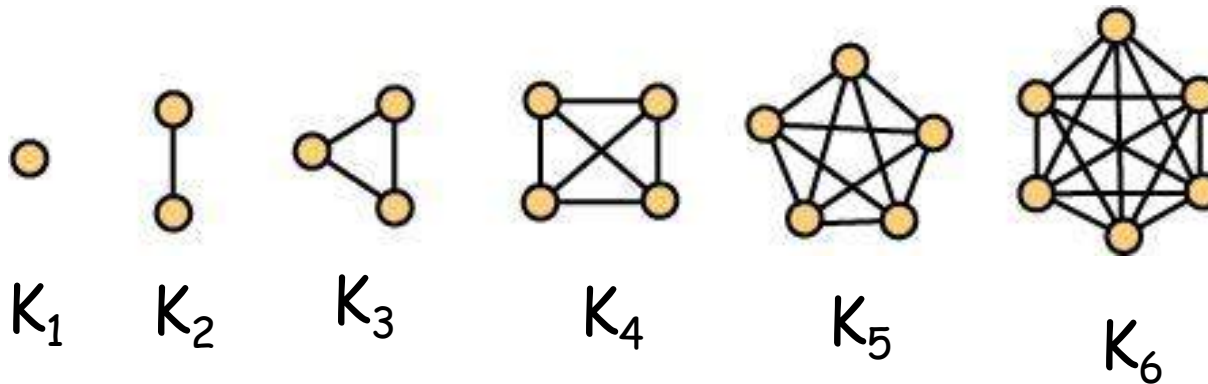
$C_5$



$C_6$

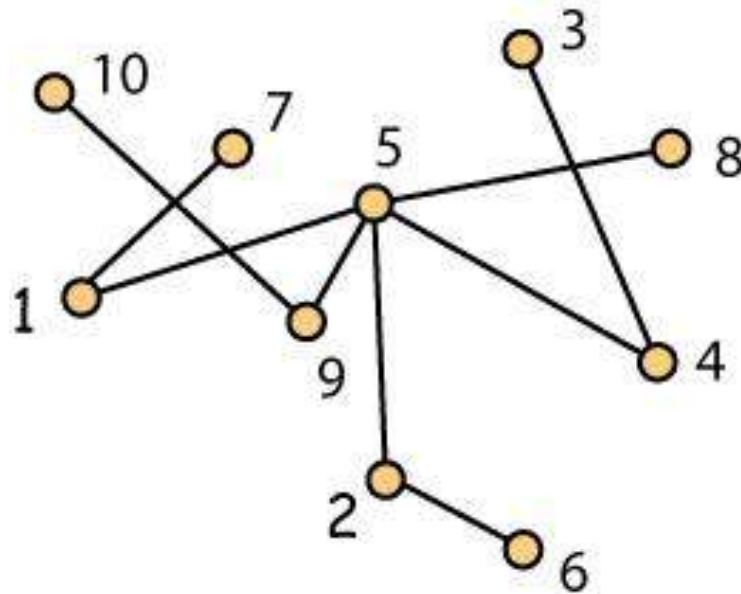
# Special Classes of Graphs (2)

**Definition** For  $n \geq 1$ ,  $K_n$  denotes a **complete graph** (also called a **clique**) on  $n$  vertices. Here are  $K_1$ ,  $K_2$ ,  $K_3$ ,  $K_4$ ,  $K_5$  and  $K_6$ .



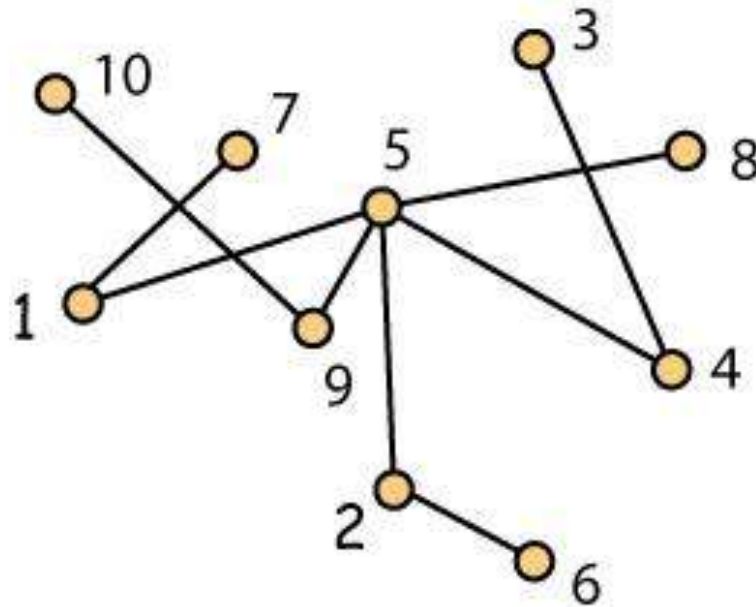
# Special Classes of Graphs (3)

**Definition** A graph  $G$  on  $n$  vertices is a **tree** if  $G$  is connected and contains no cycles.



# Properties of Trees

**Definition** When  $T$  is a tree, a vertex of degree 1 is called a leaf. This tree has five leaves: 3, 6, 7, 8, 10. The other vertices are cut vertices.



# Properties of Trees (2)

**Theorem** When  $T$  is a tree on  $n$  vertices and  $n \geq 2$ , then  $T$  has at least two leaves.

**Proof** Induction on  $n$ . Obviously true when  $n = 2$ . Assume valid when  $T$  is a tree with at least 2 and at most  $k$  vertices for some  $k \geq 2$ . Then let  $T$  be a tree with  $k + 1$  vertices. Then  $k + 1 \geq 3$ , so if  $T$  does not have 3 leaves, it has a cut vertex  $x$ . It follows that if  $C$  is a component of  $G - x$ , then  $C + x$  is a tree and has at least 2 leaves. One of these is distinct from  $x$  and is therefore a leaf in  $T$ .

# Properties of Trees (3)

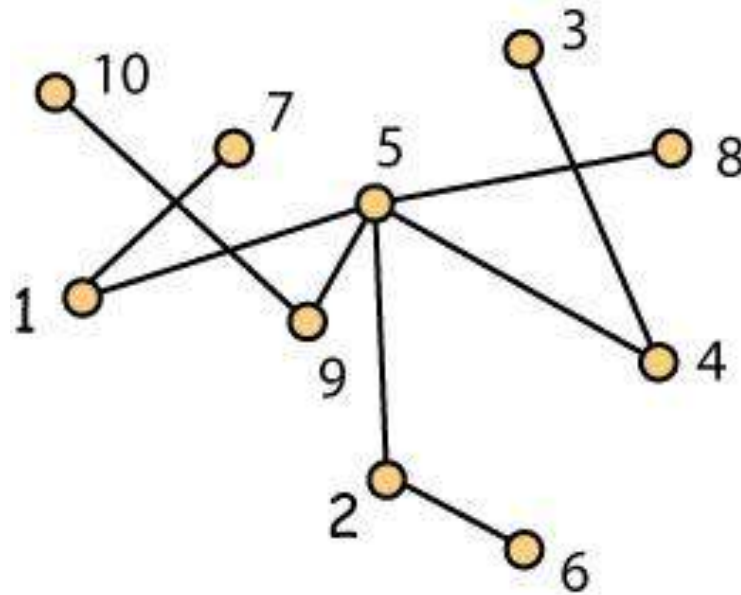
**Theorem** When  $T$  is a tree on  $n$  vertices,  $T$  has  $n - 1$  edges.

**Proof** Induction on  $n$ . True when  $n = 1$ . Now assume valid when  $n = k$  for some integer  $k \geq 1$ . Then let  $T$  be a tree on  $k + 1$  vertices. Choose a leaf  $x$  (there are at least two from which to choose). Then  $\deg(x) = 1$  while the tree  $T - x$  has  $k$  vertices and  $k - 1$  edges. Therefore  $T$  has  $(k - 1) + 1 = k$  edges.



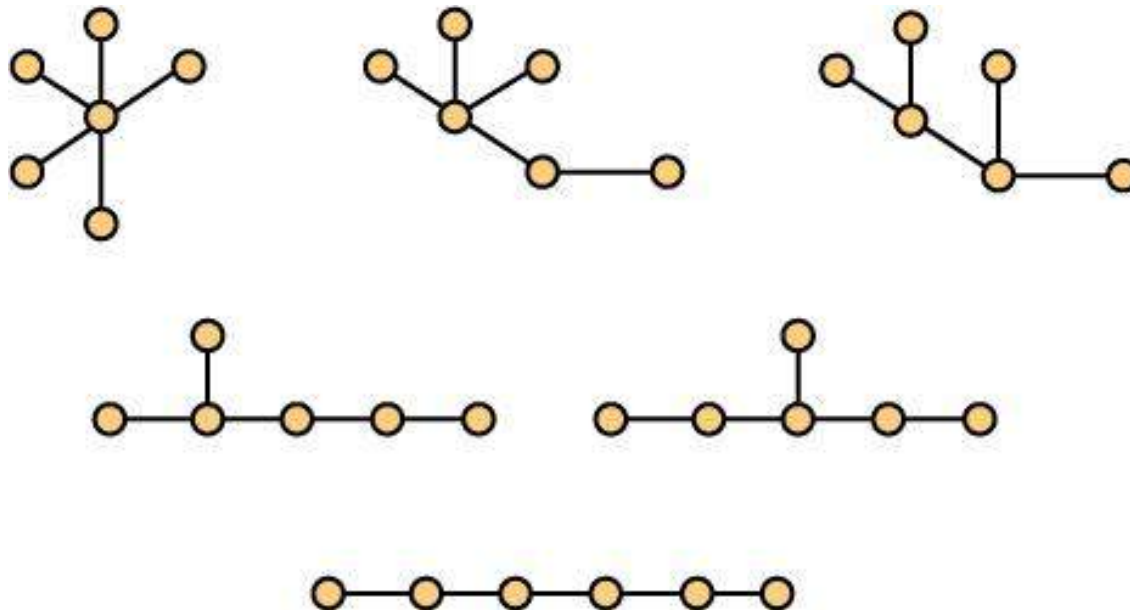
# Paths and Trees

**Theorem** When  $T$  is a tree,  $T$  is a path unless it has more than two leaves.



# Counting Trees

**Exercise** Explain why there are 6 unlabelled trees on 6 vertices. They are shown below.



# The Unlabelled Trees on 6 Vertices

**Exercise** Show that when  $1 \leq n \leq 6$ , the number of trees with vertex set  $\{1, 2, \dots, n\}$  is  $n^{n-2}$ . Actually, we did the work when  $1 \leq n \leq 5$  in class, so all you really have to do is the case  $n = 6$ .

**Remark** Later in the course, we will show that this is true for all  $n \geq 1$ .