

Analysis of Algorithm
(Solved MCS's)

CS502-Design and

LECTURE FROM (23 to 45)

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- 1. Which method is preferable for dealing with chain matrix multiplication?
- Divide and conquer strategy
- Dynamic programming formulation
- Graph theory
- Greedy Approach
- 2. Huffman algorithm produces the......prefix code tree.
- Better
- Optimal
- Worst
- Best
- 3. A...w is adjacent to vertex v if there is an edge from v to w.
  - Acyclic
  - Vertex
  - Loop
  - Cycle
- 4. Using ASCII standard the string "greedy" will be encoded with
- i. 44 bits
- ii. 120 bits
- iii. 40 bits
- iv. 48 bits
  - 5. Find the maximum value of the items which can carry using knapsack weight capacity =50

# Item weight Value

10	70
20	20
30	80
70	200

- i. 90
- ii. 280
- iii. 200
- iv. 100
  - 6. In activity scheduling algorithm, each activity is represented by a
- i. Rectangle
- ii. Square

iii. Circle iv. Triangle 7. Those problems in which greedy finds good, but not always best is called a greedy..... Heuristic ii. Solution Result iii. iv. Algorithm 8. The knapsack problem belongs to the domain of.....Problems i. Searching ii. Sorting iii. Linear solution iv. **Optimization** 9. The general coin change problem can be solved using i. Recursion ii. Greedy algorithm iii. **Dynamic programming** Divide and conquer iv. 10. Huffman algorithm produces the ...... prefix code tree V. 7 **Optimal** ii. Better iii. Best iv. Worst 11. Huffman algorithm generates an optimum......code i. **Postfix** ii. Infix iii. None of the given options iv. **Prefix** 12 ......ways of representing graphics i. 1 ii. iii. 3 iv. 13. Knapsack word originates from .....language German ii. English iii. French iv. Norwegian 14. Graphs are important......model for many application problems **Mathematical** i. ii. Unpredictable iii. Haphazard iv. Unsystematic

15. Which type of algorithm is harder to prove the correctness?

i.

Dynamic

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. Greedy	
. Divide and conquer . Brute force	
16Items are not allowed in 0/1 knapsack problem  Fractional	
. 0	
. 1	
. 0/1	
17. Matrix multiplication is a(n) operation	
Nether commutative nor associative	
. Transitive	
. Commutative	
. Associative	
18. For a Diagraph $G=(V.E)$ , Sum of in-degree $(v)$	
Not equal to sum of out-degree(v)	
- Not equal to sum of out-degree(v)	k.
• = sum of out-degree(v) $pg#115$	7
Y A A A A A A A A A A A A A A A A A A A	1
<ul><li>&lt; sum of out-degree(v)</li></ul>	Or.
	V.
<ul><li>&gt; sum of out-degree(v)</li></ul>	
19. DFS or BFS yields aof the graph.	
• Traversed tree	
Traversed tree	
• Spanning tree pg#125	
• Simple tree	
• Free tree	
20. Using ASCII code, each character is represented by a fixed-length code of bits j	per

21. In Knapsack Problem, the goal is to put items in the Knapsack

such that the value of the items is-----

ii. iii. iv.

i. ii. iii. iv.

**17.** i. ii. iii. iv.

10

character.

pg#100

-----subject to weight limit of the Knapsack.

- Minimized
- Decreased
- Maximized pg#109
- None of the above given
- **22.** Consider the string"abcdaacac",if the string is coded with ASCII codes, the message length would be---
  - 70 bits
  - 60 bits
  - 90 bits
  - 72 bits
- 23. A graph is said to be acyclic if it contains -----
  - At least one cycle
  - Exactly one cycle
  - Always more than one cycle
  - No cycles pg#116
- **24.** The number of edges that come out of a vertex is called the----- of that vertex in the digraph.
  - Post-degree
  - in-degree
  - out-degree pg#114
  - pre-degree
- **25.** If Matrix-A has dimensions "3x2" and Matrix-B has dimensions "2x3", then multiplication of Matrix-A and Matrix-B will result a new Matrix-C having dimensions.
  - 3x2
  - 2x3
  - 2x2
  - 3x3pg#87
- **26.** A/an----- is one in which you want to find, not just a solution, but the best solution.

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Optimization problem
Divide and Conquer
NP complete problem
Best problem  7. Fractional Knapsack is founded onmethod.
<ul> <li>Greedy page no109</li> <li>Recursive</li> <li>Divide and Conquer</li> </ul>
Dynamic programming
28. If the graph is represented using an adjacency matrix, then Breadth-first search takestime.
• O(E+1)
• O(V^2)
• O(V)
• O(E)
29. In inductive approach of Knapsack problem, we consider 2 cases,or
Median, Mode
• Recursive, Iterative
• Leave object, Take object pg#93
<ul> <li>Recursive, Iterative</li> <li>Leave object, Take object pg#93</li> <li>Sequentially, Parallel</li> </ul>
<b>30.</b> A Greedy algorithm can NOT be used to solve all the problems.

Dynamic programming (Google)

Memorization programming

Edit-distance programming

Storing value programming

- **31.** In Huffman encoding, the -----is the number of occurrences of a character divided by the total characters in the message.
  - Counting
  - Parsing
  - Relative Probability pg#100
  - Weight
- **32.** The Binary Tree constructed by a Huffman Encoding is a:
  - Full Binary Tree pg#102
  - Partial Binary Tree
  - Incomplete Binary Tree
  - None of the given option
- 33. Following is not the application of Edit Distance Problem.
  - Speech recognition pg#76
  - Spelling correction
  - Ascending order
  - Computational Molecular Biology
- **34.** Consider three Matrices X, Y, Z of dimensions 1x2, 2x3, 3x4 respectively. The number of multiplication of (XYZ) is:
  - 18
  - 32
  - 24
  - 30
- **35.** In ----- Knapsack Problem, limitation is that an item can either be put in the bag or not. Fractional items are not allowed.
  - a. 0
  - b. 1
  - c. 0/1 pg#91
  - d. Fractional

- **36.** An in-place sorting algorithm is one that -----uses additional array for storage.
  - Always
  - Permanently
  - Does not pg#54
  - Sometime
- **37.** If Matrix-A has dimensions "pxq" and Matrix-B has dimensions "qxr", then multiplication of Matrix-A and Matrix-B will result a new Matrix-C having dimensions.
  - P x q
  - Pxr page#84
  - qxr
  - q x p
- 38. Counting sort is suitable to sort the elements in range 1 to K.
  - K is large
  - K is small pg#71
  - K may be large or small
  - None
- 39. When matrix A of 5x 3 is multiply with matrix B of 3x 4 then the multiplication required is:

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- 15
- 12
- 36
- 60
- 40 ----- is a linear time sorting algorithm.
  - Merge sort
  - · Quick sort

- Bubble sort
- Radix sort
- 41. In Dynamic Programming approach, we do not store the solution to each sub problem in case if it reappears.
  - True
  - False page#75
- 42. Dynamic Programming approach is usually useful in solving optimization problem.
  - True page#97
- False
  - 43. Which of the following algorithm provides an optimal solution for the activity selection problem?
    - Divide and Conquer
    - Brute force
    - Greedy pg#105
    - Recursive
  - 44. A graph is -----if every vertex can reach every other vertex.
    - Connected pg#116
    - Cycle
    - Acyclic
    - Loop
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- 45. In a Huffman encoding when a new node is created by combining two nodes, the new node is placed in the \_\_\_\_\_.
- Priority queue pg#100
- Linked list
- Min heap tree
- Graph traversal
- 46. Huffman algorithm produces the \_\_\_\_\_ prefix code
- Optimal pg#105

Best
Worst
Better 47. Inalgorithm, you hope that by choosing a local optimum at each step, you will up a global
optimum.
Simple Divide and conquer  Greedy pg#97
Brute Force 48. The string "Imncde" is coded with ASCII code, the message length would bebits.
• 24
• 36
• 48
60 49. For graph traversal, breadth-first search strategy
• Is always recursive
• Cannot br recursive
Cannot be non-recursive
• Can be both recursive and non-recursive page 119
<ul> <li>50. In activity scheduling algorithm, the width of a rectangle</li> <li>Is always ignored</li> <li>Directs towards recursion</li> </ul>
• Is always ignored
Directs towards recursion
Should be maximized
• Indicates the duration of an activity pg#106
51. If the graph is represented using an adjacency list, then Breadth-first search takes
• O(V^2)

- O(V+E) pg#138
- O(E+1)
- 52. Suppose you are given infinite coins of 1,2,3, and 4. Select the ways of the minimum number CH INSTANCES of coins that required to achieve a sum of 6:
- Conceptual

- 53. sing ASCII standard the string "greedy" will be encoded with
- 48 bits Conceptual
- 120 bits
- 44 bits
- 40 bits
- 54. The Huffman codes provide a method of -----data efficiency.
- Reading/Writing
- Encoding/Decoding pg#99
- Divide/Conquer
- Inserting/Deleting
- 55. In the context of activity selection algorithm, time s dominated by sorting of the activities by-
- Start Times
- Finish Times pg#106
- **Average Times**

•	CPU Burst Times
56.	Time complexity of the "0-1" knapsack algorithm depends on
•	Number of items Capacity of the knapsack
•	Size of the Table
•	Number of items and capacity of knapsack (confirm)
57.	The greedy approach gives us an optimal solution when the coins are all powers of adenomination
4	Fixed pg#98
100	Variable
1	Constant
•	Static
58.	In Activity Selection, we say that two activities are non-interfering if their start-finish intervaloverlap
•	Do
•	Do not pg#105
•	Sometimes (3)3()4-1659294
•	Once
59.	How many steps are involved to design the dynamic programming strategy?
•	2 Vulmshelp.
•	3
•	1
•	4 pg#92
60.	Bag is a
•	type of algorithm pg#119

data structure

- program
- compiler
  - 61. If a problem is in NP-complete, it must also be in NP.
  - page#170
  - ► False
  - 62. The Huffman algorithm finds a optimal solution.
  - True
  - ➤ False
  - INSTITUTE 63. The Huffman algorithm finds an exponential solution
  - ➤ True
  - ► False pg#105
  - 64. The Huffman algorithm finds a polynomial solution
  - ➤ True google
  - ▶ False
  - 65. The greedy part of the Huffman encoding algorithm is to first find two nodes with smallest frequency.
  - ► True pg#100
  - ➤ False
  - 66. The code word assigned to characters by the Huffman algorithm have the property that no code word is the prefix of any other.
  - True pg#101
  - **►** False
  - 67. Huffman algorithm uses a greedy approach to generate a postfix code T that minimizes the expected length B (T) of the encoded string.
  - ► True
  - ► False pg#102
  - 68. Dijkestra's single source shortest path algorithm works if all edges weights are non-negative and there are negative cost cycles.
  - ► True
  - ► False pg#154

69	. The term "coloring" came form the original application which was in architectural design.
<b>&gt;</b>	True
<b>&gt;</b>	False pg#176
70	. In the clique cover problem, for two vertices to be in the same group, they must be adjacent each other.
	► True pg#176
	► False
71	. Dijkstra's algorithm is operates by maintaining a subset of vertices
	True pg#155
	► False
72	. We do sorting to,
<b>&gt;</b>	keep elements in random positions
<b>&gt;</b>	keep the algorithm run in linear order
<b>&gt;</b>	keep the algorithm run in (log n) order
	keep elements in increasing or decreasing order pg#40  After partitioning array in Quick sort, pivot is placed in a position such that
73	
73 •	. After partitioning array in Quick sort, pivot is placed in a position such that
73 •	. After partitioning array in Quick sort, pivot is placed in a position such that  Values smaller than pivot are on left and larger than pivot are on right  Values larger than pivot are on left and smaller than pivot are on right
73. <b>&gt;</b>	After partitioning array in Quick sort, pivot is placed in a position such that  Values smaller than pivot are on left and larger than pivot are on right  Values larger than pivot are on left and smaller than pivot are on right
73 <b>&gt; &gt; &gt; &gt;</b>	After partitioning array in Quick sort, pivot is placed in a position such that  Values smaller than pivot are on left and larger than pivot are on right  Values larger than pivot are on left and smaller than pivot are on right  Pivot is the first element of array
73 <b>&gt; &gt; &gt; &gt;</b>	After partitioning array in Quick sort, pivot is placed in a position such that  Values smaller than pivot are on left and larger than pivot are on right  Values larger than pivot are on left and smaller than pivot are on right  Pivot is the first element of array  Pivot is the last element of array
73 <b>&gt; &gt; &gt; &gt;</b>	After partitioning array in Quick sort, pivot is placed in a position such that  Values smaller than pivot are on left and larger than pivot are on right  Values larger than pivot are on left and smaller than pivot are on right  Pivot is the first element of array  Pivot is the last element of array  Merge sort is stable sort, but not an in-place algorithm
73	After partitioning array in Quick sort, pivot is placed in a position such that  Values smaller than pivot are on left and larger than pivot are on right  Values larger than pivot are on left and smaller than pivot are on right  Pivot is the first element of array  Pivot is the last element of array  Merge sort is stable sort, but not an in-place algorithm  True pg#54  False
73	After partitioning array in Quick sort, pivot is placed in a position such that  Values smaller than pivot are on left and larger than pivot are on right pg#48  Values larger than pivot are on left and smaller than pivot are on right  Pivot is the first element of array  Pivot is the last element of array  Merge sort is stable sort, but not an in-place algorithm  True pg#54  False  A p × q matrix A can be multiplied with a q × r matrix B. The result will be a p × r matrix C
73  74  75	After partitioning array in Quick sort, pivot is placed in a position such that  Values smaller than pivot are on left and larger than pivot are on right pg#48  Values larger than pivot are on left and smaller than pivot are on right  Pivot is the first element of array  Pivot is the last element of array  Merge sort is stable sort, but not an in-place algorithm  True pg#54  False  A p × q matrix A can be multiplied with a q × r matrix B. The result will be a p × r matrix C. There are (p · r) total entries in C and each takes to compute.
73	After partitioning array in Quick sort, pivot is placed in a position such that  Values smaller than pivot are on left and larger than pivot are on right  Values larger than pivot are on left and smaller than pivot are on right  Pivot is the first element of array  Pivot is the last element of array  Merge sort is stable sort, but not an in-place algorithm  True pg#54  False  A p × q matrix A can be multiplied with a q × r matrix B. The result will be a p × r matrix C There are (p ⋅ r) total entries in C and each takes to compute.  O (q) pg#84
73	After partitioning array in Quick sort, pivot is placed in a position such that  Values smaller than pivot are on left and larger than pivot are on right pg#48  Values larger than pivot are on left and smaller than pivot are on right  Pivot is the first element of array  Pivot is the last element of array  Merge sort is stable sort, but not an in-place algorithm  True pg#54  False  A p × q matrix A can be multiplied with a q × r matrix B. The result will be a p × r matrix C There are (p ⋅ r) total entries in C and each takes to compute.  O(q) pg#84  O(1)
73	After partitioning array in Quick sort, pivot is placed in a position such that  Values smaller than pivot are on left and larger than pivot are on right  Values larger than pivot are on left and smaller than pivot are on right  Pivot is the first element of array  Pivot is the last element of array  Merge sort is stable sort, but not an in-place algorithm  True pg#54  False  A p × q matrix A can be multiplied with a q × r matrix B. The result will be a p × r matrix C. There are (p ⋅ r) total entries in C and each takes to compute.  O (q) pg#84  O (1)  O (n²)

• Pointers pg#40

- constants
- variables
- functions

77. Merge sort requires extra array storage,

- True pg#54
- False

78. The Huffman codes provide a method of encoding data **inefficiently** when coded using

## ASCII standard.

- True
- False pg#99
- 79. Using ASCII standard the string abacdaacac will be encoded with bits
- 80 pg#99
- 160
- 320
- 100
- 80. Using ASCII standard the string abacdaacac will be encoded with 160 bits.
- True
- False pg#99
- 81. Using ASCII standard the string abacdaacac will be encoded with 10 bytes.
- True
- **♦ False pg#99**
- 82. The greedy part of the Huffman encoding algorithm is to first find two nodes with

# character frequency

- True
- \* False
- 83. Huffman algorithm uses a greedy approach to generate an prefix code T that minimizes

the expected length B (T) of the encoded string.

- **True** pg#100
- False
- 84. An optimization problem is one in which you want to find,
  - ▶ Not a solution
  - ► An algorithm
  - ► Good solution
  - ► The best solution pg#97
- 85. Although it requires more complicated data structures, Prim's algorithm for a minimum spanning tree is better than Kruskal's when the graph has a large number of vert
  - **►** True pg#149
    - ► False

AL-JUNAID TECH INSTITUTE  86. If a problem is in NP, it must also be in P.
► True
► False
► unknown pg#173
87. What is generally true of Adjacency List and Adjacency Matrix representations of graphs?
► Lists require less space than matrices but take longer to find the weight of an edge (v1,v2) pg#116
Lists require less space than matrices and they are faster to find the weight of an edge (v1,v2)
Lists require more space than matrices and they take longer to find the weight of an edge
(v1,v2)
Lists require more space than matrices but are faster to find the weight of an edge (v1,
88. If a graph has v vertices and e edges then to obtain a spanning tree we have to delete
▶ v edges.
► v – e + 5 edges
▶ v + e edges.
<b>▶</b> None of these
89. Maximum number of vertices in a Directed Graph may be  V <sup>2</sup>
► True
► False pg#115
90. The Huffman algorithm finds a (n)solution.
► Optimal pg#105 (3 (4 – 1 6 5 9 ) 9 4
► Non-optimal
► Exponential
<ul> <li>Non-optimal</li> <li>Exponential</li> <li>Polynomial</li> <li>O1. The Huffman elgorithm finds on exponential solution</li> </ul>
91. The Huffman algorithm finds an exponential solution
► True
► False pg#115

92. Edge (u, v) is a forward edge if

▶ u is a proper descendant of v in the tree

▶ v is a proper descendant of u in the tree

pg#129

► None of these

93. After partitioning array in Quick sort, pivot is placed in a position such that

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	Values smaller than pivot are on left and larger than pivot are on right (Page 48)
<b>&gt;</b>	Values larger than pivot are on left and smaller than pivot are on right
<b>&gt;</b>	Pivot is the first element of array
<b>&gt;</b>	Pivot is the last element of array
94.	Merge sort is stable sort, but not an in-place algorithm
<b></b>	True Page #54
<b>&gt;</b>	False
95.	In counting sort, once we know the ranks, we simplynumbers to their final positions in an output array.
<b>&gt;</b>	Delete
<b></b>	copy Page# 57
<b>&gt;</b>	Mark
<b>&gt;</b>	arrange
96.	. Dynamic programming algorithms need to store the results of intermediate sub-problems.
<b></b>	True pg#75
<b>&gt;</b>	False
97.	. A p $\times$ q matrix A can be multiplied with a q $\times$ r matrix B. The result will be a p $\times$ r matrix C. There are (p . r) total entries in C and each takes
<b></b>	O (q) pg#48
<b>&gt;</b>	O(1)
<b>&gt;</b>	$O(n^2)$
<b>&gt;</b>	$O(n^3)$
98.	is a graphical representation of an algorithm
<b>&gt;</b>	$\Sigma$ notation
<b>&gt;</b>	Θ notation
<b></b>	Flowchart

99. . Which of the following is calculated with  ${\bf big}~{\bf o}~{\bf notation}?$ 

► Asymptotic notation

- ► Lower bounds
- **►** Upper bounds pg#25
- ▶ Both upper and lower bound
- ► Medium bounds

100. Merge sort makes two recursive calls. Which statement is true after these recursive calls finish, but before the merge step?

- ► The array elements form a heap
- ► Elements in each half of the array are sorted amongst themselves
- ► Elements in the first half of the array are less than or equal to elements in the second half of the
- ► None of the above

What is the solution to the recurrence T(n) = T(n/2) + n, T(1) = 1101.

- ►O(logn)
- ightharpoonup O(n)

pg#37

- ►O(nlogn)
- $\triangleright$  O(2n)

102.

Consider the following Huffman Tree

The binary code for the string **TEA** is

- **► 10 00 010**
- **►**011 00 010
- ▶10 00 110
- ▶11 10 110

.A greedy algorithm does not work in phases. 103.

- ► True
- **►** False pg#9'

Can an adjacency matrix for a directed graph ever not be square in shape? 104.

- ► Yes
- **No** pg#116

105. 44. One of the clever aspects of heaps is that they can be stored in arrays without using any\_

- ► Pointers pg#40
- **▶** constants
- **▶** variables
- **▶** functions

106. .Merge sort requires extra array storage,

A	L-JUNAID TECH INSTITUTE
<b>►</b> True	Page 54
► False	
107.	.Non-optimal or greedy algorithm for money change takes
►O(k) Pg	# <b>99</b>
► O(kN)	
► O(2k)	
►O(N)	
108.	The Huffman codes provide a method of encoding data inefficiently when coded
usin	g ASCII standard.
True	
► False	Pg# 99
109.	Using ASCII standard the string abacdaacac will be encoded withbits.
1	
<b>▶80</b> pg	g# 99
▶160	
►320 ►100	
►100 110.	Using ASCII standard the string abacdaacac will be encoded with 160 bits.
110.	Oshig ASCII standard the string abacdaacac will be encoded with 100 bits.
► True	
► False (Palse)	<mark>g# 99)</mark>
111.	Using ASCII standard the string abacdaacac will be encoded with 320 bits.
► Tours	
► True ► False (P	o# 99)
112.	.Using ASCII standard the string abacdaacac will be encoded with 100 bits.
1121	veeling the est statistical and stating accordance with section with 100 estat.
➤ True	
► False (Pa	
113.	Using ASCII standard the string abacdaacac will be encoded with 32 bytes
	W
► True	g# 99)  The greedy part of the Hyffman area ding elemithm is to first find two nodes with
► False (Palse (Palse)	g# 99) 8bytes
114.	The greedy part of the Huffman encoding argorithm is to first find two nodes with
sma	<b>llest</b> frequency.
►True (Pg	<u>z# 100)</u>
► False	<del>7</del>
115.	The greedy part of the Huffman encoding algorithm is to first find two nodes with

- **♦** ►True
- **❖ False (Pg# 100)**

character frequency

116. Huffman algorithm uses a greedy approach to generate an **antefix** code T that minimizes the expected length B (T) of the encoded string.

**	$\Gamma_{r112}$
**	ı rue

- **❖ False (Pg# 102)** 
  - 117. Depth first search is shortest path algorithm that works on un-weighted graphs.
  - ► True
  - ► **False (Pg# 15**9)
    - 118. Floyd-Warshall algorithm is a dynamic programming algorithm; the genius of the algorithm is in the clever recursive formulation of the shortest path problem.
  - ► True (Pg# 162)
  - ► Flase
    - 119. .Floyd-Warshall algorithm, as in the case with DP algorithms, we avoid recursive evaluation by generating a table for
- **⋄** ▶ k
- ❖ ►True
- ❖ ►Flase
  - 120. The term coloring came from the original application which was in map drawing.
- **❖ ►**True (Pg# 176)
- **❖** ► False
  - In the clique cover problem, for two vertices to be in the same group, they must be each other.
- **♦** ► Apart from
- ♦ Far from
- ♦ Near to
- **♦ Adjacent to (Pg# 176)** 
  - 122. Fixed-length codes may not be efficient from the perspective of \_\_\_\_\_\_the total quantity of data.

# **Select correct option:**

- ► Minimizing Pg# 99
- ► Averaging
- ► Maximizing
- **▶** Summing
  - 123. In greedy algorithm, at each phase, you take the \_\_\_\_\_\_\_you can get right now, without regard for future consequences.

► Worst	
► Minimum	
<b>▶</b> Good	
<b>▶</b> Best	Page #97
124. algori	The difference between Prim's algorithm and Dijkstra's algorithm is that Dijkstra's thm uses a same key.
➤ True ➤ <b>False</b> 125.	Page# 156  If a problem is in NP-complete, it must also be in NP.
► True ► False	Page# 178
126. ►2 ►n	If there are n items, there arepossible combinations of the items.
► 2^n	Page# 92
► 3^n 127.	.Using ASCII code, each character is represented by a fixed-length code word ofbits per character.
▶4	
►6 ►8	020g#99-1659294
<b>▶</b> 10	WW. COII.
128. of the	.In Knapsack Problem, the thief's goal is to put items in the bag such that theitems does not exceed the limit of the bag.
► <mark>Value</mark>	Page #91
► Weight	
► Length	
► Balance	
129.	The knapsack problem does not belong to the domain of optimization problems.

	Page# 91
130. (rela	.In Huffman encoding, for a given message string, the frequency of occurrence ative probability) of each character in the message is determined last.
➤ True	
<b>False</b>	Page 100
131. cha	Fixed-length codes are known for easy break up of a string into its individual racters.
➤ <mark>True</mark> ➤ False	Page# 99
132.	InKnapsack Problem, limitation is that an item can either be put in the bag
not-	fractional items are not allowed.
<b>►</b> 0	
(A)	
<b>0/1</b>	
U/ I	Page# 91
Fraction:	
➤ Fraction:	al
Fractions 133.	In Knapsack Problem, value and weight both are to be under consideration.
➤ Fraction:	al
Fractions 133.	In Knapsack Problem, value and weight both are to be under consideration.
Fractions 133.  True	In Knapsack Problem, value and weight both are to be under consideration.  page #91
Fractions 133.  True  False 134.	In Knapsack Problem, value and weight both are to be under consideration.  page #91
Fractions 133.  True  False 134.  mat	In Knapsack Problem, value and weight both are to be under consideration.  page #91  Time complexity of DP based algorithm for computing the minimum cost of chain rix Multiplication is
Fractions 133.  True False 134. mat	In Knapsack Problem, value and weight both are to be under consideration.  page #91  Time complexity of DP based algorithm for computing the minimum cost of chain rix Multiplication is
Fractions 133.  True  False 134.  mat	In Knapsack Problem, value and weight both are to be under consideration.  page #91  Time complexity of DP based algorithm for computing the minimum cost of chain rix Multiplication is
Fractions 133.  True False 134. mat	In Knapsack Problem, value and weight both are to be under consideration.  page #91  Time complexity of DP based algorithm for computing the minimum cost of chain rix Multiplication is
Fractions 133.  True  False 134.  mat  log n	In Knapsack Problem, value and weight both are to be under consideration.  page #91  Time complexity of DP based algorithm for computing the minimum cost of chain
Fractions 133.  True False 134. mat log n  n	In Knapsack Problem, value and weight both are to be under consideration.  page #91  Time complexity of DP based algorithm for computing the minimum cost of chair rix Multiplication is  Page #90
Fractions 133.  True  False 134.  mat  log n  n  n2  n3	In Knapsack Problem, value and weight both are to be under consideration.  page #91  Time complexity of DP based algorithm for computing the minimum cost of chain rix Multiplication is  Page #90  In DP based solution of knapsack problem, to compute entries of V we will imply aapproach.

136.	A greedy algorithm sometimes works well for optimization problems.
≻ <mark>True Pag</mark>	ge# 9 <mark>7</mark>
► False	
137. messa	In Huffman encoding, frequency of each character can be determined by parsing t ge andhow many times each character (or symbol) appears.
Printing Incrementi	ng (Page 100)
<b>Counting</b>	(Page 100)
• Deleting	
138.	Greedy algorithm can do very poorly for some problems.
True	Page# 97
False	
139.	The Huffman codes provide a method of data efficiently.
Reading	
- Encoding	Page# 99
➤ Decoding	
➤ Printing	
140.	ke object.
140. Or Tai	ke object.
140. Or Tai	ke object.
140. Or Tai	rogramming  Page #93  Those problems in which Greedy finds good, but not always best is called a
Or Tail  Brute force  Dynamic p  141.	rogramming  Page #93  Those problems in which Greedy finds good, but not always best is called a
140. Or Tai ➤ Brute force ➤ Dynamic p 141. greedy	ke object.  Page #93  Those problems in which Greedy finds good, but not always best is called a

142. In brute force based solution of knapsack problem, we consider 2 cases, Leave object Or Take object. ► TRUE **► FALSE** Page# 97 problem, we want to find the best solution. 143. **▶** Minimization ► Averaging ► Optimization Page# 97 **►** Maximization Using ASCII standard the string abacdaacac will be encoded with 10 bytes. **►** True Page #101 ► False algorithm, you hope that by choosing a local optimum at each step, you 145. will end up at a global optimum. **►** Simple ► Non Greedy **►** Greedy Page# 97 ▶ Brute force 146. Huffman algorithm uses a greedy approach to generate an prefix code T that minimizes the expected length B (T) of the encoded string. **►** True Page# 102 ► False 147. Counting Money problem is an example which cannot be optimally solved by greedy algorithm. 1.COI Page# 97 **►** True ► False 148. Huffman algorithm generates an optimum prefix code. Page# 102 **►** True ► False

If the string "Imncde" is coded with ASCII code, the message length would be

149.

bits.

<b>▶</b> 36
► 48 6*8=48 page #99
<b>▶</b> 60
150. There arenested loops in DP based algorithm for computing the minimum cost of chain matrix multiplication.
<b>▶</b> 2
►3 Page# 90 ►4
151A number of lectures are to be given in a single lecture hall. Optimum scheduling for
this is an example of Activity selection.
► True Page# 105
► False
152. The activity scheduling is a simple scheduling problem for which the greedy algorithm approach provides a/ansolution.
► Simple
► Sub optimal
► <mark>Optimal Page# 105</mark>
▶ Non optimal
153. The string  xyz , if coded with ASCII code, the message length would be 24 bits.
► True (3*8=24) page#99
► False
154. An application problem is one in which you want to find, not just a solution, but thesolution.
► Simple
► Good Page #113
▶ Best
Suppose that a graph $G = (V,E)$ is implemented using adjacency lists. What is the complexity of a breadth-first traversal of $G$ ?

► O(|V |^2)

	1
--	---

Breadth first search is shortest path algorithm that works on un-weighted graphs  Depth first search is shortest path algorithm that works on un-weighted graphs.  Both of above are true.  Both of above are true.  I.57Forward edge is: (u, v) where u is a proper descendent of v in the tree. (u, v) where v is a proper descendent of u in the tree. (u, v) where v is a proper ancesstor of u in the tree. (u, v) where u is a proper ancesstor of v in the tree.  If you find yourself in maze the better traversal approach will be:  BFS google  BFS and DFS both are valid  Level order  DFS  In digraph G=(V,E); G has cycle if and only if  The DFS forest has forward edge.  The DFS forest has both back and forward edge  BFS forest has both back and forward edge  BFS forest has forward edge	epth first search is shortest path algorithm that works on un-weighted graphs epth first search is shortest path algorithm that works on un-weighted graphs. oth of above are true.  57Forward edge is: u, v) where u is a proper descendent of v in the tree. u, v) where v is a proper descendent of u in the tree. u, v) where v is a proper ancesstor of u in the tree. u, v) where u is a proper ancesstor of v in the tree. 58.	Breadth first  Both of above  Both of above  Ione of above  Is7F  (u, v) where  (u, v) where  (u, v) where  (u, v) where  Is8. I  BFS  BFS and DFS	arch is shortest path algorithms shortest path algorithms are true.  The are true.  The arc true are true.  The arc true are true are true.  The arc true are true are true.  The arc true arc true are true.  The arc true arc true arc true arc true.  The arc true arc true arc true arc true.  The arc true arc true arc true arc true arc true.  The arc true	nt of v in the tree of v in the tree	ee. tree.		graphs.
Depth first search is shortest path algorithm that works on un-weighted graphs.  Both of above are true.  Bone of above are true.  I.57. Forward edge is:  (u, v) where u is a proper descendent of v in the tree.  (u, v) where v is a proper descendent of u in the tree.  (u, v) where v is a proper ancesstor of u in the tree.  (u, v) where u is a proper ancesstor of v in the tree.  (u, v) where u is a proper ancesstor of v in the tree.  If you find yourself in maze the better traversal approach will be:  BFS google  BFS and DFS both are valid  Level order  DFS  In digraph G=(V,E); G has cycle if and only if  The DFS forest has forward edge.  Che DFS forest has back edge  Page#131  The DFS forest has both back and forward edge	epth first search is shortest path algorithm that works on un-weighted graphs.  oth of above are true.  one of above are true.  57Forward edge is:  u, v) where u is a proper descendent of v in the tree.  u, v) where v is a proper descendent of u in the tree.  u, v) where v is a proper ancesstor of u in the tree.  u, v) where u is a proper ancesstor of v in the tree.  58.	Depth first sees of above the	arch is shortest path algorithms are true.  The are true.  The are true.  The are true are true are true.  The are true are true are true are true.  The are true are tr	gorithm that we not of v in the tree of v in the tree	ee. tree.		graphs.
Soth of above are true.  None of above are true.  157	oth of above are true.  57	Both of above None of above 157F (u, v) where (u, v) where (u, v) where (u, v) where 158	are true.  e are true.  orward edge is:  i is a proper descendent  v is a proper descendent  i is a proper ancesstor  i is a proper ancesstor  you find yourself in r	nt of v in the tr lent of u in the of u in the tree of v in the tree	ee. tree.	weighted	STI
(u, v) where u is a proper descendent of v in the tree.  (u, v) where v is a proper descendent of u in the tree.  (u, v) where v is a proper ancesstor of u in the tree.  (u, v) where u is a proper ancesstor of v in the tree.  (u, v) where u is a proper ancesstor of v in the tree.  158. If you find yourself in maze the better traversal approach will be:  BFS google  BFS and DFS both are valid  Level order  DFS  159. In digraph G=(V,E); G has cycle if and only if  The DFS forest has forward edge.  Che DFS forest has back edge  Page#131  The DFS forest has both back and forward edge	u, v) where u is a proper descendent of v in the tree.  Page #129  u, v) where v is a proper ancesstor of u in the tree.  u, v) where u is a proper ancesstor of v in the tree.  y, v) where u is a proper ancesstor of v in the tree.  If you find yourself in maze the better traversal approach will be:  BFS google  FS and DFS both are valid  evel order  FS  JIn digraph G=(V,E); G has cycle if and only if  the DFS forest has forward edge.  the DFS forest has back edge  Page#131  the DFS forest has both back and forward edge  FS forest has forward edge  FS forest has forward edge	(u, v) where (u, v) where (u, v) where (u, v) where 158. I BFS BFS and DFS	v is a proper descendent v is a proper descendent v is a proper ancesstor v is a proper ances v is	lent of u in the of u in the tree of v in the tree	tree.		Page #
(u, v) where v is a proper ancesstor of u in the tree.  (u, v) where u is a proper ancesstor of v in the tree.  If you find yourself in maze the better traversal approach will be:  BFS google  BFS and DFS both are valid  Level order  DFS  In digraph G=(V,E); G has cycle if and only if  The DFS forest has forward edge.  Che DFS forest has back edge  Page#131  The DFS forest has both back and forward edge	u, v) where v is a proper ancesstor of u in the tree.  u, v) where u is a proper ancesstor of v in the tree.  58. If you find yourself in maze the better traversal approach will be:  BFS google  FS and DFS both are valid  evel order  FS  59. In digraph G=(V,E); G has cycle if and only if  the DFS forest has forward edge.  the DFS forest has back edge Page#131  the DFS forest has both back and forward edge  FS forest has forward edge  FS forest has forward edge	(u, v) where (u, v) where 158. I BFS BFS and DFS	is a proper ancesstor is a proper ancesstor you find yourself in r	of u in the tree			Page #
(u, v) where u is a proper ancesstor of v in the tree.  158. If you find yourself in maze the better traversal approach will be:  BFS google  BFS and DFS both are valid  Level order  DFS  159. In digraph G=(V,E); G has cycle if and only if  The DFS forest has forward edge.  Che DFS forest has back edge  Page#131  The DFS forest has both back and forward edge	If you find yourself in maze the better traversal approach will be:  BFS google  FS and DFS both are valid evel order  FS  59. In digraph G=(V,E); G has cycle if and only if the DFS forest has forward edge.  the DFS forest has back edge Page#131 the DFS forest has both back and forward edge  FS forest has forward edge  FS forest has forward edge	(u, v) where 158. I BFS BFS and DFS	i is a proper ancesstor	of v in the tree			
If you find yourself in maze the better traversal approach will be:  BFS google  BFS and DFS both are valid  Level order  DFS  159. In digraph G=(V,E); G has cycle if and only if  The DFS forest has forward edge.  Che DFS forest has back edge  Page#131  The DFS forest has both back and forward edge	If you find yourself in maze the better traversal approach will be:  BFS google  FS and DFS both are valid  evel order  FS  59. In digraph G=(V,E); G has cycle if and only if  the DFS forest has forward edge.  the DFS forest has back edge Page#131  the DFS forest has both back and forward edge  FS forest has forward edge	BFS BFS and DFS	you find yourself in r				
BFS google  BFS and DFS both are valid  Level order  DFS  159. In digraph G=(V,E); G has cycle if and only if  The DFS forest has forward edge.  Che DFS forest has back edge Page#131  The DFS forest has both back and forward edge	FS and DFS both are valid evel order FS 59. In digraph G=(V,E); G has cycle if and only if the DFS forest has forward edge.  the DFS forest has back edge Page#131 the DFS forest has both back and forward edge FS forest has forward edge	BFS BFS and DFS		naze the better			
BFS and DFS both are valid  Level order  DFS  159. In digraph G=(V,E); G has cycle if and only if  The DFS forest has forward edge.  Che DFS forest has back edge  Page#131  The DFS forest has both back and forward edge	FS and DFS both are valid evel order  FS  59. In digraph G=(V,E); G has cycle if and only if the DFS forest has forward edge.  the DFS forest has back edge  Page#131  the DFS forest has both back and forward edge  FS forest has forward edge	BFS and DFS	google		ıraversal ap	proach v	will be:
DFS  159. In digraph G=(V,E); G has cycle if and only if The DFS forest has forward edge.  Che DFS forest has back edge  Page#131  The DFS forest has both back and forward edge	evel order  FS  59. In digraph G=(V,E); G has cycle if and only if the DFS forest has forward edge.  the DFS forest has back edge  Page#131  the DFS forest has both back and forward edge  FS forest has forward edge						
In digraph G=(V,E); G has cycle if and only if The DFS forest has forward edge.  The DFS forest has back edge  Page#131  The DFS forest has both back and forward edge	59. In digraph G=(V,E); G has cycle if and only if  the DFS forest has forward edge.  the DFS forest has back edge Page#131  the DFS forest has both back and forward edge  FS forest has forward edge	)EC	both are valid				
The DFS forest has forward edge.  The DFS forest has back edge  Page#131  The DFS forest has both back and forward edge	he DFS forest has forward edge.  he DFS forest has back edge Page#131  he DFS forest has both back and forward edge  FS forest has forward edge		n dioranh G=(V F) ·G	has cycle if an	d only if		
	FS forest has forward edge						0111
	FS forest has forward edge		V.V.V	orward edge	Page#131	P.	0
	60. Back edge is:	BFS forest ha	s forward edge				

► (u, v) where v is an predcessor of u in the tree.

None of al	lbove	
161.	Cross edge is:	
➤ (u, v) wh	here u and v are not ancestor of one another	
► (u, v) wh	here u is ancesstor of v and v is not descendent of u.	
(u, v) wh	here u and v are not ancestor or descendent of one another	Page# 129
(u, v) wł	where u and v are either ancestor or descendent of one another.	
162. tree e	Kruskal's algorithm (choose best non-cycle edge) is better that edge) when the graph has relatively few edges.  Crue google	n Prim's (choose best
► Fa		1177
163. MST	What algorithm technique is used in the implementation of Kr	ruskal solution for the
<b>▶</b> Gi	Greedy Technique Page# 142	
<b>▶</b> Di	Divide-and-Conquer Technique	
► Dy	Oynamic Programming Technique	
► Th	The algorithm combines more than one of the above techniques	
164. algor	What is the time complexity to extract a vertex from the priori rithm?	ty queue in Prim's
<b>▶</b> O	O(log E)	
<b>▶</b> ( <b>∨</b>	V)	
`	V+E)	11
165.	The relationship between number of back edges and number of	f cycles in DFS is
	Both are equal	r cycles in B1 5 is,
	Back edges are half of cycles	
	Back edges are one quarter of cycles	
		Dogo# 121
166.	You have an adjacency list for G, what is the time complexity spose G^T.?	Page# 131  to compute Graph

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► (V E) ► (V) ► (V^2) 167. There is relationship between number of back edges and number of cycles in DFS ▶ Both are equal. ► Cycles are half of back edges. ► Cycles are one fourth of back edges. ► There is no relationship between back edges and number of cycles. A digraph is strongly connected under what condition? 168. ► A digraph is strongly connected if for every pair of vertices u, v e V, u can reach v ► A digraph is strongly connected if for every pair of vertices u, v e V, u can reach v and vice versa. Page# 135 ► A digraph is strongly connected if for at least one pair of vertex u, v e V, u can reach v and vice A digraph is strongly connected if at least one third pair of vertices u, v e V, u can reach v and vice versa. 169. In in-place sorting algorithm is one that uses arrays for storage: ► An additional array ► No additional array Page#54 ▶ Both of above may be true according to algorithm ► More than 3 arrays of one dimension. In stable sorting algorithm 170. ▶ One array is used ► In which duplicating elements are not handled. ► More than one arrays are required. ▶ Duplicating elements remain in same relative position after sorting. 171. Which sorting algorithm is faster: ► O(n^2) ► O(nlogn)

versa.

**▶** O(n+k)

172.	In Quick sort algorithm, constants hidden in T(n lg n) are
1/2.	in Quick sort argorithm, constants indden in 1(ii ig ii) are
► La	arge
► Medium	
► Not know	n
► Small  173. Quick pivot element	sort is based on divide and conquer paradigm; we divide the problem on base of and:
➤ There is e ➤ No work	xplicit combine process as well to conquer the solution. s needed to combine the sub-arrays, the array is already sorted he sub arrays
► None of a	
174 Dijkst	ra's algorithm:
	y approach to find all shortest paths greedy and Dynamic approach to find all shortest paths
► Has both  ► Has gree	
► Has both	greedy and Dynamic approach to find all shortest paths
► Has both  ► Has gree  154  ► Has both	greedy and Dynamic approach to find all shortest paths
► Has both  ► Has gree  154  ► Has both vertices.	greedy and Dynamic approach to find all shortest paths dy approach to compute single source shortest paths to all other vertices  greedy and dynamic approach to compute single source shortest paths to all other
► Has both  ► Has gree  154  ► Has both vertices.	greedy and Dynamic approach to find all shortest paths  dy approach to compute single source shortest paths to all other vertices  Page 19 19 19 19 19 19 19 19 19 19 19 19 19
► Has both  ► Has gree  154  ► Has both vertices.	greedy and Dynamic approach to find all shortest paths dy approach to compute single source shortest paths to all other vertices  greedy and dynamic approach to compute single source shortest paths to all other may be stable sort:
► Has both : ► Has gree  154 ► Has both vertices. 175. Which	greedy and Dynamic approach to find all shortest paths dy approach to compute single source shortest paths to all other vertices  greedy and dynamic approach to compute single source shortest paths to all other may be stable sort:
► Has both green  Has green  154  ► Has both vertices.  175. Which  ► Bubble so  ► Insertion  ► Both of a	greedy and Dynamic approach to find all shortest paths dy approach to compute single source shortest paths to all other vertices  greedy and dynamic approach to compute single source shortest paths to all other may be stable sort:  ort sort bove  page# 54
► Has both : ► Has gree  154 ► Has both vertices. 175. Which ► Bubble so ► Insertion	greedy and Dynamic approach to find all shortest paths dy approach to compute single source shortest paths to all other vertices  greedy and dynamic approach to compute single source shortest paths to all other may be stable sort:  ort sort bove  page# 54
► Has both  Has gree  154  ► Has both vertices. 175. Which  ► Bubble so ► Insertion  ► Both of a ► Selection	greedy and Dynamic approach to find all shortest paths dy approach to compute single source shortest paths to all other vertices  greedy and dynamic approach to compute single source shortest paths to all other may be stable sort:  ort sort bove  page# 54 sort
<ul> <li>► Has both</li> <li>► Has green</li> <li>154</li> <li>► Has both vertices.</li> <li>175. Which</li> <li>► Bubble so</li> <li>► Insertion</li> <li>► Both of a</li> <li>► Selection</li> <li>176. In the a</li> </ul>	greedy and Dynamic approach to find all shortest paths dy approach to compute single source shortest paths to all other vertices  greedy and dynamic approach to compute single source shortest paths to all other may be stable sort:  ort sort bove  page# 54 sort
► Has both  ► Has gree  154  ► Has both vertices. 175. Which  ► Bubble so ► Insertion  ► Both of a ► Selection  176. In the a phase; we ge	greedy and Dynamic approach to find all shortest paths dy approach to compute single source shortest paths to all other vertices  greedy and dynamic approach to compute single source shortest paths to all other may be stable sort:  ort  sort  bove
<ul> <li>► Has both</li> <li>► Has green</li> <li>154</li> <li>► Has both vertices.</li> <li>175. Which</li> <li>► Bubble so</li> <li>► Insertion</li> <li>► Both of a</li> <li>► Selection</li> <li>176. In the a</li> </ul>	greedy and Dynamic approach to find all shortest paths dy approach to compute single source shortest paths to all other vertices  greedy and dynamic approach to compute single source shortest paths to all other may be stable sort:  ort sort bove page# 54 sort  analysis of Selection algorithm, we eliminate a constant fraction of the array with each et the convergent
► Has both  ► Has gree  154  ► Has both vertices. 175. Which  ► Bubble so ► Insertion  ► Both of a ► Selection  176. In the a phase; we go  ► linear	greedy and Dynamic approach to find all shortest paths  dy approach to compute single source shortest paths to all other vertices  greedy and dynamic approach to compute single source shortest paths to all other  may be stable sort:  ort  sort  bove
► Has both   ► Has green  154  ► Has both vertices.  175. Which ► Bubble so ► Insertion   ► Both of a ► Selection  176. In the a phase; we go ► linear ► arithmetic	greedy and Dynamic approach to find all shortest paths  dy approach to compute single source shortest paths to all other vertices  greedy and dynamic approach to compute single source shortest paths to all other  may be stable sort:  ort  sort  bove
► Has both  ► Has gree  154  ► Has both vertices. 175. Which  ► Bubble so ► Insertion  ► Both of a ► Selection  176. In the a phase; we go  ► linear  ► arithmetic ► geometri ► exponent	greedy and Dynamic approach to find all shortest paths  dy approach to compute single source shortest paths to all other vertices  greedy and dynamic approach to compute single source shortest paths to all other  may be stable sort:  ort  sort  bove

178. Dijkstra's algorithm is operates by maintaining a subset of vertices

True

Page# 155

False

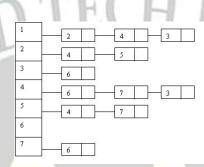
179. The difference between Prim's algorithm and Dijkstra's algorithm is that Dijkstra's algorithm uses a different key.

**►** True

Page# 156

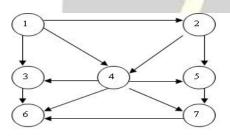
► False

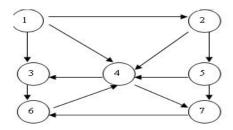
180. Consider the following adjacency list:

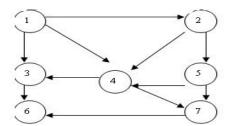


STITUM

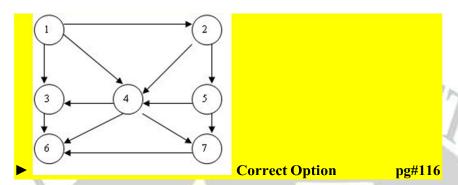
Which of the following graph(s) describe(s) the above adjacency list?







304-1659294



- 181. We do sorting to,
  - ▶ keep elements in random positions
  - ► keep the algorithm run in linear order
  - ♦ keep the algorithm run in (log n) order
  - ★ keep elements in increasing or decreasing order

Page# 40

- 182. After partitioning array in Quick sort, pivot is placed in a position such that
  - ❖ Values smaller than pivot are on left and larger than pivot are on right Page# 48
  - ❖ ► Values larger than pivot are on left and smaller than pivot are on right
  - ❖ ► Pivot is the first element of array
  - ❖ ► Pivot is the last element of array
- 183. . Merge sort is stable sort, but not an in-place algorithm
- ❖ ► True

(Page# 54)

- ♦ False
- 184. In counting sort, once we know the ranks, we simply \_\_\_\_\_numbers to their final positions in an output array.
  - **❖** ▶ Delete
  - **⇔ ►** copy

(Page# 57)

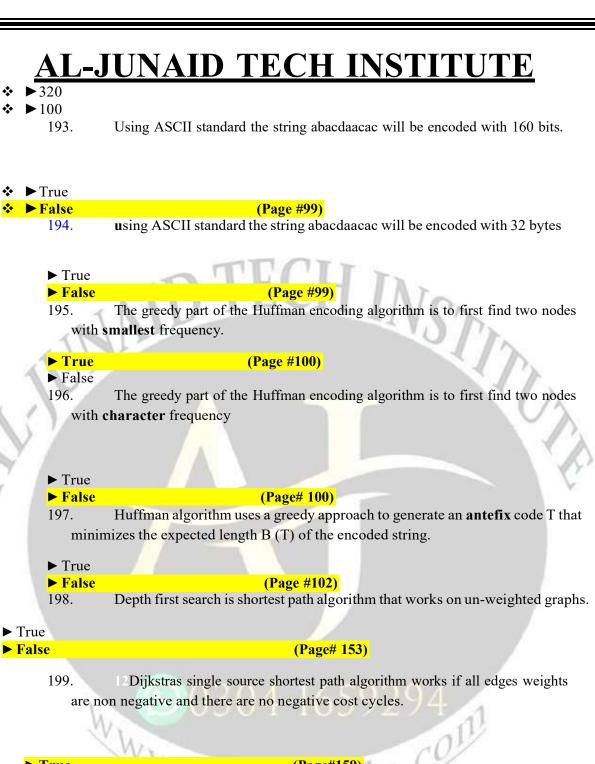
- ♦ Mark
- ❖ ▶ arrange
- 185. Dynamic programming algorithms need to store the results of intermediate sub-problems.
  - **♥** Iru

(Page# 75)

➤ False

	(Da za# 94)	
	<b>►</b> O (q) (Page# 84)	
	$ ightharpoonup O(n^2)$	
*	$ ightharpoonup O(n^3)$	
18	7 is a graphical representation of an algorithm	
<b>*</b>	$ ightharpoonup \Sigma$ notation	
	▶ <sup>1</sup> notation	
•	► Flowchart	
•	► Asymptotic notation	
	Which of the following is calculated with <b>big o notation</b> ?	
<b>&gt;</b>	► Both upper and lower bound  ► Medium bounds  189. Merge sort makes two recursive calls. Which statement is true	after these
٠.	recursive calls finish, but before the merge step?	
	► The array elements form a heap	
*	► The array elements form a heap Elements in each half of the array are sorted amongst themselves	ne second half of
<ul><li>*</li><li>*</li></ul>	<ul> <li>▶ The array elements form a heap</li> <li>▶ Elements in each half of the array are sorted amongst themselves</li> <li>▶ Elements in the first half of the array are less than or equal to elements in the the array</li> <li>▶ None of the above</li> </ul>	ne second half o
<ul><li>*</li><li>*</li></ul>	<ul> <li>▶ The array elements form a heap</li> <li>▶ Elements in each half of the array are sorted amongst themselves</li> <li>▶ Elements in the first half of the array are less than or equal to elements in the array</li> </ul>	ne second half of
<b>* * *</b>	<ul> <li>▶ The array elements form a heap</li> <li>▶ Elements in each half of the array are sorted amongst themselves</li> <li>▶ Elements in the first half of the array are less than or equal to elements in the the array</li> <li>▶ None of the above</li> </ul>	ne second half o
<b>* * *</b>	<ul> <li>► The array elements form a heap</li> <li>► Elements in each half of the array are sorted amongst themselves</li> <li>► Elements in the first half of the array are less than or equal to elements in the the array</li> <li>► None of the above 190. Non-optimal or greedy algorithm for money change takes</li> <li>► O(k) (Page #99)</li> <li>► O(kN)</li> </ul>	ne second half o
	<ul> <li>► The array elements form a heap</li> <li>► Elements in each half of the array are sorted amongst themselves</li> <li>► Elements in the first half of the array are less than or equal to elements in the array</li> <li>► None of the above <ul> <li>190. Non-optimal or greedy algorithm for money change takes</li> </ul> </li> <li>► O(k)  (Page #99) <ul> <li>► O(kN)</li> <li>► O(2k)</li> </ul> </li> </ul>	ne second half o
	<ul> <li>► The array elements form a heap</li> <li>► Elements in each half of the array are sorted amongst themselves</li> <li>► Elements in the first half of the array are less than or equal to elements in the array</li> <li>► None of the above <ul> <li>190. Non-optimal or greedy algorithm for money change takes</li> </ul> </li> <li>► O(k) (Page #99)</li> <li>► O(kN)</li> <li>► O(2k)</li> <li>► O(N)</li> </ul>	
<ul><li>*</li><li>*</li></ul>	<ul> <li>► The array elements form a heap</li> <li>► Elements in each half of the array are sorted amongst themselves</li> <li>► Elements in the first half of the array are less than or equal to elements in the the array</li> <li>► None of the above 190. Non-optimal or greedy algorithm for money change takes</li> <li>► O(k) (Page #99)</li> <li>► O(kN)</li> <li>► O(2k)</li> <li>► O(N)</li> <li>191. The Huffman codes provide a method of encoding data inefficence in the array are less than or equal to elements in the themselves</li> </ul>	
	<ul> <li>► The array elements form a heap</li> <li>► Elements in each half of the array are sorted amongst themselves</li> <li>► Elements in the first half of the array are less than or equal to elements in the array</li> <li>► None of the above <ul> <li>190. Non-optimal or greedy algorithm for money change takes</li> </ul> </li> <li>► O(k) (Page #99)</li> <li>► O(kN)</li> <li>► O(2k)</li> <li>► O(N)</li> <li>191. The Huffman codes provide a method of encoding data ineffice coded using ASCII standard.</li> </ul>	
	<ul> <li>► The array elements form a heap</li> <li>► Elements in each half of the array are sorted amongst themselves</li> <li>► Elements in the first half of the array are less than or equal to elements in the array</li> <li>► None of the above 190. Non-optimal or greedy algorithm for money change takes</li> <li>► O(k) (Page #99)</li> <li>► O(kN)</li> <li>► O(2k)</li> <li>► O(N)</li> <li>191. The Huffman codes provide a method of encoding data ineffice coded using ASCII standard.</li> <li>► True</li> </ul>	
	<ul> <li>► The array elements form a heap</li> <li>► Elements in each half of the array are sorted amongst themselves</li> <li>► Elements in the first half of the array are less than or equal to elements in the array</li> <li>► None of the above <ul> <li>190. Non-optimal or greedy algorithm for money change takes</li> </ul> </li> <li>► O(k) (Page #99)</li> <li>► O(kN)</li> <li>► O(2k)</li> <li>► O(N)</li> <li>191. The Huffman codes provide a method of encoding data ineffice coded using ASCII standard.</li> </ul>	iently when

**►80 ►**160 (Page #99)



► True (Page#159)

► False

200. Dijkestra s single source shortest path algorithm works if all edges weights are negative and there are no negative cost cycles.

► True

► False (Page #159)

201. Floyd-Warshall algorithm is a dynamic programming algorithm; the genius of the algorithm is in the clever recursive formulation of the shortest path problem.

	202.	Floyd-Warshall algorithm, as in the case with DP algorithms, we avoid rsive evaluation by generating a table for
		isive evaluation by generating a table for
	►k	
	► True	
	► False	
	203.	the term coloring came from the original application which was in map
	draw	ving.
	<b>►</b> True	(Page #176)
	► False	
	204.	In the clique cover problem, for two vertices to be in the same group, they the each other.
	mus	t deeach other.
1		Apart from
JI 3		Far from
1		➤ Near to  Adjacent to (Page# 176)
N	205.	Fixed-length codes may not be efficient from the perspective of
DA.	£	the total quantity of data.
Salaat aar	most ontic	
Select cor	тесі орис	
► Minimi	zing	(Page #99)
► Averag	ing	
► Maxim	izing	
<b>▶</b> Summi	ng	
, Summin		
	206.	In greedy algorithm, at each phase, you take theyou can get right
	now	, without regard for future consequences.
	1	0304-1037274
► Worst	4	
► worst		1 - 00
► Minimu	ım	VIII mahal Di
► Good		www.vulmshelp.com
<b>▶</b> Best		(Page #97)
	207.	The difference between Prim's algorithm and Dijkstra's algorithm is that
		stra s algorithm uses a same key.
	► True	( Do zo#15 ()
	<b>► False</b> 208.	( Page# 156)  If a problem is in NP-complete, it must also be in NP.
		1

► Tru	e	(Page #178)
► Fals	re.	
Tais		
	209.	If there are n items, there arepossible combinations of the items.
*	▶2	
*	▶n	
<b>*</b>	<b>▶</b> 2^n	(Page# 92)
*	▶3^n	· · · · · · · · · · · · · · · · · · ·
	210.	Using ASCII code, each character is represented by a fixed-length code word
	of	bits per character.
*	<b>▶</b> 4	TO THE HIM.
*	▶6	VIII THOU
*	▶8	(Page #99)
*	<b>▶</b> 10	(t age #22)
•	211.	In Knapsack Problem, the thief's goal is to put items in the bag such that the
	4111	
	Nalus	_of the items does not exceed the limit of the bag.
	► Value	(Page# 91)
. \	► Weight	
*	Length	
V **	Balance	
Y	212.	The knapsack problem does not belong to the domain of optimization
	proble	ems.
*	<b>►</b> True	
**	<b>►</b> False	(Page# 91)
	<b>213.</b>	In Huffman encoding, for a given message string, the frequency of occurrence
	(relativ	ve probability) of each character in the message is determined last.
*	<b>►</b> True	
<b>*</b>	<b>►</b> False	(Page #100)
	214.	Fixed-length codes are known for easy break up of a string into its individual
	charac	
	1	1000304-1659794
<b>►</b> Tru	<mark>e</mark>	(Page# 99)
► False	. //	W W 2 2011
1 415		12 6
	215.	InKnapsack Problem, limitation is that an item can either be put in the
		not-fractional items are not allowed.
*	<b>▶</b> 0	
*	<b>▶</b> 1	
•		
*	<b>▶</b> 0/1	(Page# 91)
*	► Fractional	(* "ge" > 1)
•	216.	The term "coloring" came from the original application which was in
		ectural design.
*	► True	
<b>*</b>	► False	(Page# 173)
	217	In Vacance le Duchlam, reduce and revisible hothers to be under consideration

# **AL-JUNAID TECH INSTITUTE** (page# 91)

*	
	►False
	218. Time complexity of DP based algorithm for computing the minimum cost of chain matrix Multiplication is
*	•
*	
*	n2
	▶n3 (Page# 90)
	219. In DP based solution of knapsack problem, to compute entries of V we will
	imply a/an approach.
*	► Subjective
*	
	▶Brute force
	► Combination
	A greedy algorithm sometimes works well for optimization problems.
	► True (Page# 97)
	▶ False
	221. In Huffman encoding, frequency of each character can be determined by
	parsing the message andhow many times each character (or symbol)
1	appears.
S	app man
7.	Y'
٠.	N. Dallacian
<b>*</b>	► Printing
*	► Incrementing
.*.	Counting (Depot 100)
<b>*</b>	( was tri)
<u> </u>	<b>▶</b> Deleting
<b>.</b>	➤ Deleting 222Greedy algorithm can do very poorly for some problems.
* *	<ul> <li>▶ Deleting</li> <li>222Greedy algorithm can do very poorly for some problems.</li> <li>▶ True (Page# 97)</li> </ul>
<b>.</b>	<ul> <li>▶ Deleting     222Greedy algorithm can do very poorly for some problems.</li> <li>▶ True     (Page# 97)</li> <li>▶ False</li> </ul>
*	<ul> <li>Deleting     222Greedy algorithm can do very poorly for some problems.</li> <li>► True</li></ul>
* * *	<ul> <li>Deleting         <ul> <li>222Greedy algorithm can do very poorly for some problems.</li> </ul> </li> <li>▶ True</li></ul>
* * * * * *	<ul> <li>Deleting     222Greedy algorithm can do very poorly for some problems.</li> <li>▶ True</li></ul>
* * * * *	<ul> <li>Deleting         <ul> <li>222Greedy algorithm can do very poorly for some problems.</li> </ul> </li> <li>▶ True</li></ul>
* * * * *	<ul> <li>Deleting         <ul> <li>222. Greedy algorithm can do very poorly for some problems.</li> </ul> </li> <li>True</li></ul>
* * * * * *	<ul> <li>Deleting         <ul> <li>222Greedy algorithm can do very poorly for some problems.</li> </ul> </li> <li>True</li></ul>
* * * * *	<ul> <li>Deleting         <ul> <li>222. Greedy algorithm can do very poorly for some problems.</li> </ul> </li> <li>True</li></ul>
< <  < <  < <  < <  < <  < <  < <  < <  < <  < <  < <  < <  < <  < <  < <  < <  < <  < <  < <  < <  < <  < <  < <  < <  < <  < <  < <  < <  < <  < <  < <  < <  < <  < <  < <  < <  < <  < < <  < < <  < < < < < < < <>  < < < < < < < < < < < < < < < < < < <	<ul> <li>Deleting         <ul> <li>222Greedy algorithm can do very poorly for some problems.</li> </ul> </li> <li>True</li></ul>
* * * * * * * * * * * * * * * * * * *	<ul> <li>Deleting</li></ul>
* * * * *	<ul> <li>Deleting</li></ul>
* * * * * * * * * * * * * * * * * * *	<ul> <li>Deleting         <ul> <li>222. Greedy algorithm can do very poorly for some problems.</li> </ul> </li> <li>▶ True (Page# 97)</li> <li>▶ False         <ul> <li>223. The Huffman codes provide a method of data efficiently.</li> </ul> </li> <li>▶ Reading <ul> <li>▶ Encoding</li> <li>▶ Printing</li> <li>224. In based solution of knapsack problem, we consider 2 cases, Leave object Or Take object.</li> </ul> </li> <li>▶ Brute force <ul> <li>▶ Dynamic programming</li> <li>(Page #93)</li> </ul> </li> <li>225. Those problems in which Greedy finds good, but not always best is called a</li> </ul>
* * * * * * * * * * * * * * * * * * *	<ul> <li>Deleting</li></ul>
*	<ul> <li>Deleting</li></ul>
*	<ul> <li>Deleting</li></ul>

In brute force based solution of knapsack problem, we consider 2 cases, Leave object Or Take object. **►**TRUE **► FALSE** (Page# 97) 227. problem, we want to find the best solution. **►** Minimization ► Averaging **▶** Optimization (Page#97) **►** Maximization 228. Using ASCII standard the string abacdaacac will be encoded with 10 bytes. **►**True (Page# 101) ❖ ► False algorithm, you hope that by choosing a local optimum at each step, 229. you will end up at a global optimum. **►** Simple ► Non Greedy (Page# 97) ▶Brute force 230. Huffman algorithm uses a greedy approach to generate an prefix code T that minimizes the expected length B (T) of the encoded string. **►**True (Page #102) ► False How many elements do we eliminate in each time for the Analysis of Selection 231. algorithm? (Page #37) ▶ n / 2 elements  $\triangleright$  (n/2) + n elements ▶ n / 4 elements ▶ 2 n elements 232. Slow sorting algorithms run in,  $ightharpoonup T(n^2)$ (Page #39) **►** T(n) ► T( log n) ightharpoonup T(n log n) Counting sort is suitable to sort the elements in range 1 to k: 233. ► K is large ► K is small (Page# 57) ► K may be large or small ► None Heaps can be stored in arrays without using any pointers; this is due to the 234. nature of the binary tree, **▶** left-complete ▶ right-complete ► tree nodes ► tree leaves 235. Sieve Technique can be applied to selection problem? **►** True (Page#35)

▶ False

(Page# 40)

236. A heap is a left-complete binary tree that conforms to the

increasing order onlydecreasing order only

► heap order

▶ (log n) order 237. Divide-and-conquer as breaking the problem into a small number of **▶** pivot ► Sieve ► smaller sub problems (Page #34) **▶** Selection 238. In Sieve Technique we do not know which item is of interest **►** True (Page# 34) ► False 239. The recurrence relation of Tower of Hanoi is given below  $T(n)=\{1 \text{ if } n=1 \text{ and } n=1$ 2T(n-1) if n > 1 In order to move a tower of 5 rings from one peg to another, how many ring moves are required? ▶ 16 ▶ 10 ▶ 32 **▶** 31 240. For the heap sort, access to nodes involves simple operations. **►** arithmetic ( Page# 41) **▶** binary ► algebraic **▶** logarithmic 241. For the sieve technique we solve the problem, **►** recursively (Page# 34) ▶ mathematically **▶** precisely **▶** accurately ❖ 242. The sieve technique works in as follows (Page #34 **▶** phases **▶** numbers **▶** integers **▶** routines s a left-complete binary tree that conforms to the heap 243. A (an) order **▶** heap (Page# 40) ▶ binary tree ▶ binary search tree **▶** array 244. The sieve technique is a special case, where the number of sub problems is just ▶5 **▶** many

1	AL-JUNAID TECH INSTITUTE
	► 1 (Page #34)
	245. Analysis of Selection algorithm ends up with,
	► $T(n)$ ► $T(1/1+n)$ ► $T(n/2)$ ► $T((n/2)+n)$ (Page #37)
	For the heap sort we store the tree nodes in
×.	<ul> <li>▶ level-order traversal</li> <li>▶ in-order traversal</li> <li>▶ pre-order traversal</li> <li>▶ post-order traversal</li> <li>247. The reason for introducing Sieve Technique algorithm is that it illustrates a very important special case of,</li> </ul>
	► divide-and-conquer (Page# 34)
61	► decrease and conquer
1	▶ greedy nature
PA	► 2-dimension Maxima
A.	248. Theta asymptotic notation for T (n):
,	<ul> <li>▶ Set of functions described by: clg(n)Set of functions described by clg(n)&gt;=f(n) focls</li> <li>▶ Theta for T(n)is actually upper and worst case comp</li> <li>▶ Set of functions described by:</li> <li>▶ clg(n</li> </ul>
	249. Sieve Technique applies to problems where we are interested in finding a
	single item from a larger set of
	n items (Page# 34)
	phases phases
	pointers 0504-1659294
	➤ constant 250. Memorization is?
	230. Wellorization is:
	provious results for rutare use
	id this unnecessary repetitions by writing down the results of recursive calls and
	hem up again if we need them later (Page# 47) the the process accurate
	f the above
	251. Quick sort is

One example of in place but not stable algorithm is

(Page#57)

► Stable & in place

► Not stable but in place
► Stable but not in place

252.

► Some time stable & some times in place

► Merger Sort	
► Quick Sort	(Page# 54)
► Continuation Sort	
► Bubble Sort	
253.	. Continuation sort is suitable to sort the elements in range 1 to k
► K is Large	
► K is not known	
► K may be small or l	
► K is small	(Page# 57)
254.	Which may be a stable sort?
► Merger  • Insertion	
► Insertion ► Both above	(Daga 454)
None of the above	(Page #54)
255.	. An in place sorting algorithm is one that usesarrays for storage
► Two dimensional ar	· · · · · · · · · · · · · · · · · · ·
► More than one array	
► No Additional Arr	
None of the above	(1 age not)
256.	Continuing sort has time complexity of?
► O(n	
► O(n+k)	
► O(nlogn)	
► O(k)	
257.	single item from a larger set of
▶ n items	(Page# 34)
▶ phases	
<b>▶</b> pointers	
▶ v constant	
258.	For the Sieve Technique we take time
► T(nk)	( Page# 34)
ightharpoonup T(n/3)	1.
► n^2	WY
► n/3	William
259.	One Example of in place but not stable sort is  (Page# 54)
<b>►</b> Quick	(Page# 54)
► Heap ► Merge ► Bubble 260. Factorial (n){	Consider the following Algorithm:
if (n=1)	
return 1	

else

return (n \* Factorial(n-1))

Recurrence for the following algorithm is:

- **♦** T(n) = T(n-1) + 1
- **♦** T(n) = nT(n-1) + 1
- T(n)=T(n-1)+n
- T(n)=T(n(n-1))+1
  - 261. Due to left complete nature of binary tree, the heap can be stored in
- ❖ ► Arrays

(Page #40)

- ❖ ► Structures
- ❖ ► Link List
- ♦ Stack
  - 262. What type of instructions Random Access Machine (RAM) can execute?
- ❖ ► Algebraic and logic
- ❖ Geometric and arithmetic
- **❖** ► Arithmetic and logic

(Page# 10)

- ❖ ► Parallel and recursive
  - 263. What is the total time to heapify?
- **❖** O(log n) (Page 43)
- $\bullet$   $\triangleright$  O(n log n)
- $\bullet$   $\triangleright$  O(n<sup>2</sup> log n)
- $\bullet$   $\triangleright$   $O(\log^2 n)$ 
  - 264. word Algorithm comes from the name of the muslim author
- **♦** Abu Ja'far Mohammad ibn Musa al-Khowarizmi.
  - 265. Al-Khwarizmi's work was written in a book titled
- ❖ ▶al Kitab al-mukhatasar fi hisab al-iabr wa'l-muqabalah
  - 266. Random access machine or RAM is a/an
- ❖ ► Machine build by Al-Khwarizmi
- ❖ ► Mechanical machine
- ❖ ► Electronics machine

memory.

❖ ► Mathematical model (Page# 10)

A RAM is an idealized machine with

random-access

- **❖** ▶ 256MB
- **♦** ► 512MB
- **♦** ► an infinitely large

(Page #10)

- **♦** ► 100GB
  - 268. What will be the total number of max comparisons if we run brute-force maxima algorithm with n elements?
- **♦** ▶ *n*
- $n^{\frac{n}{2}}$
- **♦** ▶ <sup>n</sup> (Page# 14)

269. Consider the following code: For(j=1; j< n; j++)For(k=1; k<15;k++) For(l=5; l<n; l++) Do something constant(); What is the order of execution for this code. ightharpoonup O(n) $\bullet$   $\triangleright$   $O(n^3)$  $\bullet$   $\triangleright$   $O(n^2 \log n)$  $\bullet$   $\triangleright$   $O(n^2)$ 270. Is it possible to sort without making comparisons? ► Yes (Page# 57) ► No 271. When we call heapify then at each level the comparison performed takes time ightharpoonup It will take  $\Theta$  (1) (Page# 43) ❖ ► Time will vary according to the nature of input data ❖ ► It can not be predicted ❖ ► It will take  $\Theta$  (log n) 272. In Quick sort, we don't have the control over the sizes of recursive calls ► True (Page# 40) ► False ❖ ► Less information to decide ❖ ► Either true or false 273. For Chain Matrix Multiplication we cannot use divide and conquer approach because, **▶** We do not know the optimum k (Page# 86) ❖ ► We use divide and conquer for sorting only ❖ ► We can easily perform it in linear time ❖ ► Size of data is not given The Knapsack problem belongs to the domain of problems. 274. **▶** Optimization (Page# 91) ❖ ► NP Complete ❖ ► Linear Solution ❖ ► Sorting

Suppose we have three items as shown in the following table, and suppose the

275.

capacity of the knapsack is 50 i.e. W = 50.

Item	Value	Weight
1	60	10
2	100	20
3	120	30

The optimal solution is to pick

- **❖** ► Items 1 and 2
- **❖** ► Items 1 and 3
- **❖** ► Items 2 and 3
- ♦ None of these

who invented the quick sort

**♦ C.A.R. Hoare** 

277. main elements to a divide-and-conquer

**❖ Divide, conquer, combine** 

278. Mergesort is a stable algorithm but not an in-place algorithm.

- **❖ ►**True (Page# 54)
- ❖ ► false

279. Counting sort the numbers to be sorted are in the range 1 to k where k is small.

- **♦ ► True** (Page# 57)
- False

280. In selection algorithm, because we eliminate a constant fraction of the array with each phase, we get the

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- **❖** Convergent geometric series (Page 37)
- ❖ Divergent geometric series
- ❖ ► None of these

281. ... In RAM model instructions are executed

**❖** ► One after another

(Page# 10)

❖ ► Parallel

	AL-J	UNAID TECH INSTITUTE
	♦ Concurr	rent
	❖ ►Randon	1
	282.	Due to left-complete nature of binary tree, heaps can be stored in
*	► Link list	
*	► Structure	
*	► Array	(Page# 40)
*	None of abo	· · ·
•	283.	The time assumed for each basic operation to execute on RAM model of
		•
*	► Infinite	itation is
Ĭ		TO THE HITE
**	Continuous	
**	<b>►</b> Constant	(Page# 10)
**	, arraore	
	284.	If the indices passed to merge sort algorithm are not equal, the algorithm may
		immediately.
*	► True	Α Α
*	<b>►</b> False	(Page# 28)
1	285.	Brute-force algorithm uses no intelligence in pruning out decisions.
**	<b>►</b> True	(Page #18)
**	►False	
Y	286.	In analysis, the Upper Bound means the function grows asymptotically no
	faster	than its largest term.
*	<b>►</b> True	(Page #24)
*	►False	
	287.	For small values of n, any algorithm is fast enough. Running time does
	becom	e an issue when n gets large.
**	<b>►</b> True	(Page #14)
*	<b>►</b> False	
	288.	In simple brute-force algorithm, we give no thought to efficiency.
**	<b>►</b> True	(Page# 11)
*	►False	- 000 1 1 ( FOD 0 1
	289.	The ancient Roman politicians understood an important principle of good
	algorit	hm design that is plan sweep algorithm
*	► True	
**	<b>►</b> False	(Page# 27)
	290.	In 2d-space a point is said to be if it is not dominated by any other
		n that space.
*	► Member	in that option
*	► Minimal	
*	► Maximal	(Page# 11)
*	► Iviaximai  ► Joint	(Fage# 11)
**		An elecuithm is a mothematical autity that is described and are a second
	291.	An algorithm is a mathematical entity that is dependent on a specific
	progra	mming language.
	<b>N</b> T	
*	► True  False	(Page# 7)

292. The running time of an algorithm would not depend upon the optimization by the compiler but that of an implementation of the algorithm would depend on it.

**	<b>►</b> True	(Page# 13)
*	►False	
	293.	F (n) and g (n) are asymptotically equivalent. This means that they have
	ess	entially the same for large n.
<b>*</b>	► Results	-50
*	►Variable	es
*	<b>►</b> Size	, 111   FL FI   ATA
**	<b>►</b> Growth	rates (Page#23)
	294.	8n2 + 2n - 3 will eventually exceed $c2*(n)$ no matter how large we make $c2$ .
*	►True	(Page# 25)
<b>*</b>	►False	
	295.	If we associate (x, y) integers pair to cars where x is the speed of the car and y
	is t	he negation of the price. High y value for a car means acar.
*	►Fast	
*	►Slow	
*	► Expensi	ve / / /
**	<b>►</b> Cheap	(Page# 11)
	296.	The function $f(n) = n(\log n + 1)/2$ is asymptotically equivalent to n log n. Here
	Up	per Bound means the function f(n) grows asymptoticallyfaster than
	n lo	og n.
		► More
		➤ Quiet
		Not (Page# 24)
		➤ At least
	297.	Counting Money problem is an example which cannot be optimally solved by
	gre	edy algorithm.
**	<b>►</b> True	(Page# 97)
*	►False	000041007274
	298.	Huffman algorithm generates an optimum prefix code.
**	<b>►</b> True	(Page #102)
*	►False	W Mark 1 al Div
	299.	If the string "Imnode" is coded with ASCII code, the message length would be
		bits.
*	▶24	
*	▶36	
*	<b>►48</b> (6	5*8=48)
*	<b>▶</b> 60	
	300.	There arenested loops in DP based algorithm for computing the
	miı	nimum cost of chain matrix multiplication.
*	▶2	
**	▶3	(Page# 90)

**	▶5	
	301.	Inductive approach to compute entries of V is implied inbased so
	of kr	napsack problem.
*	►Brute force	e
**	<b>▶</b> Dynamic	programming (Page #93)
	302.	A number of lectures are to be given in a single lecture hall. Optimum
	sche	duling for this is an example of Activity selection.
**	<b>►</b> True	(Page# 105)
*	<b>►</b> False	
	303.	The activity scheduling is a simple scheduling problem for which the gre
	algoi	rithm approach provides a/ansolution.
*	<b>►</b> Simple	
*	►Sub optim	al
**	<b>▶</b> Optimal	(Page# 105)
*	►Non optin	
4	304.	The string  xyz , if coded with ASCII code, the message length would be
	bits.	
*	►True (3	<mark>*8=24)</mark>
*	►False	
Y	305.	An application problem is one in which you want to find, not just a solut
	but tl	hesolution.
*	<b>►</b> Simple	
*	► Good	(Page# 113)
*	▶Best	
*	►Worst	
	306.	A free tree with n vertices has exactlyedges.
	a	. n
	ь	. n+1
	c. n-1	page 142 3 0 4 - 1659294
	1	Vi.
	307.	Kruskal`s algorithm works by addingin increasing order of weig
		test edge first)
	( 8	Ulmshell
	a	. Verticals
	b. <mark>Edges</mark>	page 147
		page 147
	• Trees	
	• Weights	

a. Size

#### b. Generalization

page 135

- Optimization
- Connection

The process of updating estimates in Dijkstra's algorithm is called \_\_\_\_\_

- a. Updating
- b. Amendment

#### c. Relaxation

page 154

- Insertion
  - 310. An un-weighted graph can be considered as a graph in which every edge has weight Unit.
    - 2
    - b. 5
    - c. 3

#### d. 1

page 153

- The breadth-first-search algorithm is a shortest-path algorithm that works on graphs.
  - a. Weighted
  - b. Directed
- c. Un-weighted

page 153

d. Un-directed

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#### NAID TECH INST Which activity creates a unique cycle in a free tree: adding any vertex b. adding any sub tree c. adding root d. adding any edge page 142 313. The relationship between number of back edges and Number of cycles in DFS is, Both are equal b. Back edges are half of cycles c. Back edges are one quarter of cycles d. There is no relationship between no. of edges and cycles From given algorithms which one considered as best for finding the shortest-path: DFS Bellman-Ford algorithm Dijkstra's algorithm page 154 / google **BFS** 315. Overall time for Kruskal algorithm is $\Theta(\log E)$ b. $\Theta(E \log V)$ page 149 / google $\Theta(E \log E)$ d. $\Theta(V \log E)$ Dijkstra's Algorithm is used to solve problems. 316.

b. Single-source shortest path

**page 154** 

- Multi-source shortest path
- Sorting & searching

a. All-pair shortest path

317. A digraph is strongly connected under what condition?

- A digraph is strongly connected if for every pair of vertices u, v e V, u cancreach v.
- A digraph is strongly connected if for every pair of vertices u, v e V, u can
   reach v and vice versa.

  Page 135
- A digraph is strongly connected if for at least one pair of vertex u, v e V, u can reach v and vice versa.
- A digraph is strongly connected if at least one third pair of vertices u, v e V, u can reach v and vice versa.



318. Floyd-Warshall Algorithm is based on \_\_\_\_\_

	a.]	Dynamic Programming page 161
	•	Greedy Approach
	•	Divide and Conquer
	•	Complexity theory
a. b.	An	319. In strong components algorithm, vertices are sorted inorder of finish times.
4	c.	Decreasing page 141 / google
	d.	Strong  320is commonly the running time of Dijkstra's Algorithm using the binary heap method.
	a.	<b>Θ</b> (E log V) page 156 Θ (V log )
	•	$\Theta$ (log E)
	•	$\Theta$ (Blog V)
	9	Which technique is used in the implementation of Kruskal solution for the MST?  Greedy Technique page 142 / google
		Divide-and-Conquer Technique
		Greedy Technique page 142 / google  Divide-and-Conquer Technique  Dynamic Programming Technique
	c.	The election could be a see of the characteristics. Divide and
	d.	The algorithm combines more than one of the above techniques i.e. Divide-and-
		Conquer and Dynamic Programming
		322. A fully connected undirected graph of 5 nodes will haveedges.
	4	

• 5

10 n(n-1)/2

d. 15

323. Prim's algorithm is based on \_\_\_\_\_strategy.

a. Greedy

page 150 / google

- b. Dynamic programming
- c. Divide and Conquer
- d. Exponential

We say that two vertices u and v are mutually \_\_\_\_\_ if u can reach v and vice versa.

- a. Crossed
- b. Forward
- c. Reachable

page 135

d. Not Reachable

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325.	In Dijkstra's algorithm, initi	ally the estimated	value from so	ource vertex to any
vertex	v is:			

- a. Zero (0)
- b. One (1)
- c. Minus one (-1)
- **d.** Infinity  $(\infty)$

page 154

326. A strongly connected component only apply to:

a. Directed Graph

page 135 / google

- Undirected Graph
- Minimum Spanning Tree
- Breadth First Search

327. A graph may contain

- Exactly one MST
- No MST
- One or zero MST
- More than one MST

google

328. In \_\_\_\_algorithm(s), at any time, the subset of edges A forms a single tree.

- a. Kruskals
- b. Prim's

**page 149** 

- c. kruskal's and Prim's
- d. BFS

329. The \_\_\_\_\_ given by DFS allow us to determine a number of things about a graph or

Digraph

a. color stamps

b. time stamps	page 130

- c. line stamps
- d. node stamps

330. In Kruskal's algorithm, the vertices will be stored in\_\_\_\_\_

a. Links

b. Sets page 147

- nodes
- Loops

Keeping in mind the shortest path, if given scenarios occur in computer networks like the internet where data packets have to be routed. The vertices are \_\_\_\_\_. Edges are

which may be wired or wireless.

#### a. Routers, communication links

**page 153** 

- Internet, routers
- Communication links, routers
- Routers, internet

_	332. digraj	In computing the strongly connected components of a digraph. vertices of the
are	into	subsets
a.	Joined	
<mark>b.</mark>	Partitioned	d page 135
c.	Deleted	TO THE HIM
d.	Created	INDITION INCA
	17	
	333.	In undirected graph, by convention all the edges are callededges.
a.	Forward	
1	> 1/A	
b.	Back	page 130
y	Carren	
	c. Cross	
		Forward and back
	334.	In Timestamped DFS, No back edges means
a.	1 cycle	
<mark>b.</mark>	no cycles	page 131
	c. DFS	1000274
		In Prim's algorithm, we start with the vertex r, it can be any vertex
	d. BFS	VIII mahelP
	335.	. In Prim's algorithm, we start with thevertex r, it can be any vertex.
•	Pivot	
•	Leaf	
•	negative	
d	Root	nage 149

. The ancestor and descendent relation can be nicely inferred by the lemma.

- addition a.
- b. division
- c. parenthesis

page 129

- d. Node
  - Which of the following statement is true? 337.
- Kruskal algorithm is multiple source technique for finding MST
- Kruskal's algorithm is used to find minimum spanning tree of a graph, time complexity of this algorithm is 0(EV)
- Both I and II
- Kruskal's algorithm (choose best non-cycle edge) is better than Prim's (choose best Tree edge) when the graph has relatively few edges.

See option B of Sr 46

. The time complexity to compute Graph transposes G^T is (V+E), if you have \_ 338. for G.

a. an adjacency list

page 138

- Array list
- ©0304-1659294 Www.vulmshelp.com complete list
- stack

AL-JUNAID TECH INSTITUTE  339. In Prim's algorithm, If there is no edge from u to a vertex in S. we set the key value to
key value to
• 0
• 1
• -1
<ul> <li>d. ∞ page 151</li> <li>340 What is the time complexity to extract a vertex from the priority queue in Prim's algorithm?</li> </ul>
• O(log E)
• O(V)
• O(V+E)
d. O(logV) page 152
341 Edge weights can be interpreted as distance
2 11 2 2 2 2 2 2 2 2 2 2 2 2 2 2 2 2 2
• in breadth-First Search
<ul><li>in breadth-First Search</li><li>in Queue's</li></ul>
<ul> <li>in breadth-First Search</li> <li>in Queue's</li> <li>c. in the shortest-paths page 153</li> </ul>
<ul> <li>in breadth-First Search</li> <li>in Queue's</li> <li>c. in the shortest-paths page 153</li> <li>d. in depth-First Search</li> <li>342. Thegiven by DFS allow us to determine whether the graph contains</li> </ul>
<ul> <li>in breadth-First Search</li> <li>in Queue's</li> <li>c. in the shortest-paths page 153</li> <li>d. in depth-First Search</li> <li>342. Thegiven by DFS allow us to determine whether the graph contains</li> </ul>
<ul> <li>in breadth-First Search</li> <li>in Queue's</li> <li>in the shortest-paths page 153</li> <li>d. in depth-First Search</li> <li>342. The given by DFS allow us to determine whether the graph contains any</li> </ul>
<ul> <li>in breadth-First Search</li> <li>in Queue's</li> <li>c. in the shortest-paths page 153</li> <li>d. in depth-First Search</li> <li>342. Thegiven by DFS allow us to determine whether the graph contains</li> </ul>
<ul> <li>in breadth-First Search</li> <li>in Queue's</li> <li>c. in the shortest-paths page 153</li> <li>d. in depth-First Search</li> <li>342. The given by DFS allow us to determine whether the graph contains any</li> <li>cycles.</li> <li>a. Order</li> </ul>

Ш

343. Dijkstra's algorithm works on a weighted directed graph G(V,E) in which all weights

are non-nega	ative.
--------------	--------

a. Vertices

b. Edges page 154

- nodes
- links

344. By breaking any edge on a cycle created in free tree, the free\_\_\_\_\_is restored.

- a. Edge
- b. Tree page 142
  - c. Cycle
  - d. Vertex

345. . Bellman-Ford algorithm is used to solve\_\_\_\_\_problems.

- a. All pair shortest path
- b. Single source shortest path

google

mshelp.com

- c. Flow of networking
- d. Double source shortest path

346. From the given options which one is correct regarding the time complexity of Dijkstra's algorithm

- O(N)
- O(N^3)
- O(N^2)

#### d. O(log N)

google

347. In which algorithm, information of shortest path is propagated sequentially along each

shortest path in the graph.

#### a. Bellman-Ford

page 16

- 52) Brute-force technique
- 53) Dijkstra's
- 54) Prim's

348. In Timestamped DFS-cycles lemma, if edge (u, v) is a tree, forward or cross edge, then a. f[u] < f[v]

b. f[u] > f[v]

page 130

- 52)  $f[u] \le f[v]$
- 53)  $f[u] \ge f[v]$

349. Bellman-Ford Algorithm does not allow G(graph) to have \_\_\_\_

a. positive cost cycles

#### b. negative cost cycles

**page 159** 

- c. negative weights edges
- d. positive weights edges

350. Kruskal algorithm (choose best non cycle edge) is better than Prims (choose best tree edge) when the \_\_\_\_\_has relatively few\_\_\_\_\_

a. tree, edges

b.	graph edges	see option D of Sr 23
•	tree, branches	
•	graph, branches	
	351. Finding the faste vertices and we use	r result of the shortest path –
	a. Single-pair shorte	
	b. Two pair shortest	
b.	All pair shortest paths proble both I and II	lem pa
Y	352. There are no undirected graph.	edges in
a.	Forward	
b.	Back	
c.	Cross page 1	30
d.	Both forward and back	
	W W	
	W.V	ulmshe

in the sense that it is possible from any location in the Networks are network to

reach any other location in the digraph

a. Complete **page 135** 

- b. Incomplete
- c. Not graphs
- d. Transportation
- 354. . Dijkstra's algorithm:
- a. Has greedy approach to find all shortest paths
- FECH INSTANCE b. Has both greedy and Dynamic approach to find all shortest paths
- Has greedy approach to compute single source shortest paths to all other

vertices page 154

- e. Has both greedy and dynamic approach to compute single source shortest paths to all other vertices
- Bellman-Ford algorithm is slower than 355.
- a. Brute-force technique
- page 159 b. Dijkstra's
- 58) Prim's
- 59) Graph Algorithm
  - 356. Dijkstra's Algorithm cannot be applied on
  - a. directed and weighted graphs
- b. graphs having negative weight function
- unweighted graphs

- · undirected and unweighted graphs
  - 357. Which of the following is used in the data structure for implementing Dijkstra's Algorithm?
    - a. Max heap
    - b. Stack's
    - c. Circular queue

d. Priority queue

Google

358. In Generic approach determining of Greedy MST, we maintain a subset A of

a. Edges

**page 143** 

- Vertices
- Cycles
- Paths
  - 359. In Dijkstra's algorithm the estimated value of source vertex d[s] is

a. Equal to 0

page 154

- Equal to 1
- Greater than 0
- Greater than 1

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360. Dijkstra's algorithm is a simple\_\_\_\_\_ algorithm for computing the single-source

shortest-paths to all other vertices	shortest-	paths	to	all	other	vertices
--------------------------------------	-----------	-------	----	-----	-------	----------

- a. Greedy page 154
- c. Bellman-Ford
- d. Divide and conquer
- e. Brute-Force
  - 361. In Timestamped DFS. If there is a back edge (u, v) then v is an ancestor of u and by following tree edge from v to u, we get\_\_\_\_\_.
  - a. Nothing
  - b. a cycle page 131
    - c. a line
    - d. a graph
    - There exists a unique path between any vertices of a free tree.
  - a. One
  - b. Two page 142 ) 4-1059294
  - Three
  - Four
    - technique is look like propagating wave-front outward.
  - a. Generic traversal
  - b. Breadth first traversal page 117
    - c. Depth first traversal

d. Time stamp traversal

364. You have an adjacency list for G, what is the time complexity to compute Graph transpose G^T?

a. (V+E) page 138

- V.E
- V
- Ε

In the shortest-paths problem, we are given a weighted of G = (V, E).

page 153 a. Directed graph

- 65) Line graph
- 66) Un-directed graph
- 67) Weighted graph

366. Equivalence relation partitions the vertices into classes of mutually reachable

vertices and these are the strong components

a. Variance

... equivalence
d. Non classes b. Equivalence

367.	Overall Running time of Prim's algorithm is
	a. Θ(ElogE)
	b. $\Theta(ElogV)$
c. Θ(('	V+E)logV) page 152
d. Θ((V	+E)logE)
368.	Forgraphs, there is no distinction between forward and back edges.
a. large	
b. directed	
c. und	irected page 130
d. Medi	um
369.	For graphs, there is no distinction between forward and back edges.
a. Undi	rected page 130
• direc	eted
• sma	
• large	
	the 00304-1659294
vertices.	W. 2011
a. Set	www.vulmshelp.com
b. Graph	GIIIISIICA
c. Tree	
d. Fores	st page 147
371.	In strong components algorithm, the form of graph is used in which all theof
original gra	ph G have been reversed in direction.

a. Vertices

<mark>b.</mark>	. Edges		page 138
c.	Both edge	es & vertices	
d.	Trees		
	372.	In computing the	c
partiti	ioned into s	ubsets.	LL
a.	weakly co	nnected	
	10		
b.	strongly c	onnected	pa
c.		/ A	
d.	Worst		
	373.	Timestamp structu	ure of
	of a		
digra	ph.		
<mark>a.</mark>	DFS	google	e
1 D	Ed		
b. B			
	oth DFS &	BFS	
d. M	IST	WW	
		WW.VI	ulm

	374. not.	In Prim's algorithm, if the	ne color of a vertex is	then it is in S otherw
a.	White			
b.	Gray			
	c. Black	page 151		
	d. Blue	an TI	CHD	
	375.  a. are not use	111	communication and transp	portation networks.
71	10,			177
1	b. are used	page 135		
c.	parts are used	$\wedge$		
d.	final value is	used		
	376.	Once you enter a strong o	component, every vertex is	n the component is
	a. not reacha	ble		
	<mark>b. reachable</mark>	page 137		
c.	reachable son	ne times 030		14 11
<b>u.</b>	377.	In Kruskal's algorithm, not induce a/an	the next edge is added to v 	riable set A, if its adding
a.	Vertex			
b.	Edge			
	c. Cycle	page 147		
	d. Tree			

378. Problems such as the shortest route between cities can be solved efficiently by modelling the road map as a\_\_\_\_\_\_\_.

- a. Tree
- b. Graph page 153
- c. Linked list
- d. Stack

379. In Bellman-Ford Algorithm, relaxation applies to every edge of the graph and repeat this \_\_\_\_\_\_time.

- a. E-1
- b. E + 1
- c. V+1
  - d. V 1 page 159

380. A topological sort of a DAG is a \_\_\_\_\_\_ordering of the vertices of the DAG such

that for each edge (u, v), u appears before v in the ordering.

a. Linear page 134

- b. Parallel
- c. Sequence
- d. Non-linear

	381.	Adding any edge to a free tree creates a unique
a.	Vertex	
b.	Edge	
	c. Cycle	page 142
d. \$	Strong comp 382.	onent Back edge is:
	a. (u, v) who	ere v is an ancestor of u in the tree. Page 128
	11	
b.	(u, v) where	u is an ancestor of v in the tree.
c.	(u, v) where	v is a predecessor of u in the tree.
d.	(u, v) where	u is a mid of v in the tree.
Y	383.	In Kruskal's algorithm, the nextis not added to viable set A, if its ac
		,
ind	uce a/an cycl	
ind		
ind	a. Vertex	le.
ind	a. Vertex	le.
ind	a. Vertex  b. Edge c. Cycle	le.
ind	a. Vertex  b. Edge c. Cycle d. Tree  384.	page 147 Which of the following statement is false about Dijkstra's Algorithm?
ind	a. Vertex  b. Edge c. Cycle d. Tree  384.	page 147  00304-1659294
	<ul><li>a. Vertex</li><li>b. Edge</li><li>c. Cycle</li><li>d. Tree</li><li>384.</li><li>a. It can be</li></ul>	page 147 Which of the following statement is false about Dijkstra's Algorithm?
b.	<ul> <li>a. Vertex</li> <li>b. Edge</li> <li>c. Cycle</li> <li>d. Tree</li> <li>384.</li> <li>a. It can be</li> <li>It is used to see</li> </ul>	page 147  Which of the following statement is false about Dijkstra's Algorithm?  applied on graphs having a negative weight function  generally graphs having a negative weight function
b. c.	a. Vertex  b. Edge c. Cycle d. Tree  384.  a. It can be  It is used to so	page 147  Which of the following statement is false about Dijkstra's Algorithm?  applied on graphs having a negative weight function  go solve Single-source shortest path a weighted directed graph
b. c.	a. Vertex  b. Edge c. Cycle d. Tree  384.  a. It can be  It is used to so	page 147  Which of the following statement is false about Dijkstra's Algorithm?  applied on graphs having a negative weight function  generally graphs having a negative weight function

a.	Start	p:	age 138
b.	Finish		
c.	Both start a	& finish	
	4 4		
d.	Middle		
	386.	. The tricky part of	al
	of an	T	H)
odao v	vill arasta a	cycle in viable set A.	L JLI
euge v	viii create a t	Lycie ili viable set A.	
	11		
a.	Kruskal's		page
b.	Prim's		
c.	Both Krusk	al's and Prim's	
d.	DFS		
Y			
	387. vertice		ents are not
	vertice		
reacha	ability.		
	Strongly or	nnootod	page 139
	Strongly co		page 139
	Weakly con	inected	
	First two		
d.	Last two	W	
		WW.VU	1
		·VU	Ims

<b>AL</b>	-JUNAID TECH INSTIT
	In Prim's algorithm, we will make use of
	a. Stack
	b. Priority Queue page 150
c.	Array
d.	List
	The component digraph is necessarily
a.	straight
b.	cyclic
1	c. Acyclic page 136
D.	d. Strong
У	390. A free tree with nhave exactly n - 1
a.	vertices, nodes
b.	edges, vertices
c.	nodes, vertices

page 142

If you find yourself in maze the better traversal approach will be:

google (always return with short part)

d. vertices, edges

391.

c. Level order

392.

edge is:

Forward

(u, v) where u is proper descendent of v in the tree

b. (u, v) where v is proper descendent of u in the tree page 129

a. BFS

d. DFS

	-J U I	NAID IECH INSIIIU.
c.	(u, v) where v	is proper ancestor of u in the tree
d.	(u, v) where u	is proper ancestor of v in the tree
	393.	. cross edge is :
a.	(u, v) where u	and v are not ancestor of one another
b.		is ancestor of v and v is not descendent of u
c.	(u. v) where u	and v are not ancestor or descendent of one another page 129
d.		and v are either ancestor of descendent of one another . The running time of Bellman-Ford algorithm is
a.	$\Theta(V + E)$	
b.	$\Theta(E+E)$	
1	c. Θ(VE)	page 159
G.	d. $\Theta(V+V)$	
7	395.	In Bellman-Ford Algorithm, relaxation applies toof the graph.
	a. <mark>Every edge</mark>	page 159
b.	Every Vertices	
c.	Only First Ver	tices
d.	Only First edg	e
	396.	If a subset of edges A is visible for building MST, it cannot contain a/an
a.	Vertex	w.vulmshelp.co
b.	Edge	ww.vulmshelp.com
	c. Cycle	page 143
	d. Graph	
	397.	. In Timestamped DFS-cycles lemma, if edge (u, v) is a back edge, then

 $a. \quad f[u] \le f[v]$ 

# <u>AL-JUNAID TECH INSTITUTE</u>

b.	f[u]	> f	$[\mathbf{v}]$
----	------	-----	----------------

c.  $f[u] \leq f[v]$ 

page 130

- d.  $f[u] \ge f[v]$ 
  - 398. In digraph G=(V,E);G has cycle if and only if
  - a. The DFS forest has forward edge.

b. The DFS forest has back edge

page 131

- c. The DFS forest has both back and forward edge
- d. BFS forest has forward edge

399. The key[u] is the weight of the going from u to any vertex in S.

- a. lightest edge page 151
- b. edge
- c. lighter edge
- d. heavier edge

400. In Bellman-Ford Algorithm, path consists of at most edges.

- a. V+1
- b. V-1

page 160

- c. E+1
- d. E-1

401. According to parenthesis lemma, vertex u is a descendent of v vertex it and only if;

a. [d[u], f[u]] ⊆ [d[v], f[v]]

**page 129** 

- b.  $[d[u], f[u]] \supseteq [d[v], f[v]]$
- c. Unrelated
- d. Disjoint

	402. Edge weights can be interpreted as distance
a.	in breadth-First Search
b.	in Queue's
c.	in the shortest-paths page 153
d.	in depth-First Search
1	<ul><li>403algorithm allows negative weights edges and no negative cost cycles.</li><li>a. Brute-force technique</li></ul>
13	
) b.	Bellman-Ford page 159
c.	Dijkstra's
d.	Print
	<ul><li>404. In Huffman encoding, the characters with smallest probabilities are placed at the depth of the tree.</li><li>a. Minimum</li></ul>
	a. Minimum  b. Average
•	Maximum page 102
•	Root 0304-1659294
	405. There are ways to representing graphs.
•	405. There are ways to representing graphs.  3
•	GIIIISIICE
•	2 page 116
	a. 4 406. Each time we traverse graph by Breadth-first search algorithm, we count the distance from

a. Starting node
Neighbours of the starting node page 117
Right most node
Left most node
407. In Activity Selection problem, intuitively
There are always short activities as input
Short activities are not attractive
Duration of the activities does not matter
We do not like long activities page 105
408. The activity scheduling is a simple scheduling problem for which the gree algorithm approach provides a/ansolution.
a. Simple
b. Sub-optimal
Optimal page 105
c. Non optimal
409. The is a problem for which the greedy algorithm approach pro an optimal solution.
Activity selection page 105
Dynamic programming
Knapsack problem
NP complete problem

a. Adjacency list

b. Incidence matrix

• Adjacency list, Adjacency Matrix page 116

• No way to represent

411.

Queue Data Structure work on ----- principles

Δ

FIFO(first in first out)

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- LIFO(last in first out)
- JLO(just in last out)
- LOFI(last out first in)
  - 412. Graphs can be represented by an----- and -----
- queue, stack
- adjacency list , adjacency matrix

page 116

- a. adjacency right, adjacency left
- b. Binary, linear
- 413. Identify a TRUE statement about Knapsack.
  - a. The Knapsack Problem does not belongs to the domain of optimization problems
- The Knapsack Problem belongs to the domain of optimization problems page 91
  - i. The Knapsack Problem cannot be solved by using Dynamic programming
- The Knapsack Problem is optimally solved by using Brute force algorithm

414. Which of the following is true about graph?

- a. A graph may contain no edges and many vertices
- A graph may contain many edges and no vertices
  - 415. Which type of algorithm is harder to prove the correctness?
- Greedygoogle
- Brute force
- Dynamic programming
- Divide and Conquer
  - 416. The running time of brute-force algorithm to solve Knapsack problem is-----
- O(h1)
- O(n3)
- O(n!)
- $O(2^n) pg#92$

# 0304-1659294

- 417. If a matrix has three rows and two columns, the dimension of matrix will be:
- 3x2 Conceptual
- 2x3
- 3x3
- 2x2

418. The Probability in Huffman encoding is the number of occurrences of a character divided by the total-

----in the message.

- c. Numbers
- d. Frequencies
- e. Strings
- f. Characters pg#100
- 419. In recursive formulation of Knapsack Problem: V
  [0, j] =------for j>=0
- -1
- 0 pg#93
- 1
- 2
- 420. The Knapsack Problem belongs to the domain of----- problem.
  - g. Optimization pg#91
  - h. Sorting
  - i. Linear solution
  - j. Searching
- 421. An optimization problem is one in which you want to find the ----- solution.
  - k. Simple
  - 1. Good
  - m. Best pg#97

• Worst

422. Huffman algorithm generates an co optimum----- de.

• Prefix page 102

- Postfix
- Infix
- None of the above
  - 423. Which of the following algorithms solves the Fractional Knapsack Problem more effectively?

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- Divide and Conquer
- Greedy algorithm page 109
- Dynamic programming
- Backtracking
  - 424. One of the limitation in 0/1 Knapsack is that an item can either be -----in the bag or not.

relp.con

- Use
  - Put page 91
  - Move
- Store
  - 425. In Dynamic Programming based solution of Knapsack Problem, if we decide to take an object'i' then we gain ----------------.

- a. W(Total weight of Knapsack)
- b. V(Total value of all items)
- vi(value of object i) page 93
- None of the given option
  - 426. An activity scheduling algorithm, the width of a rectangle --------
    - n. Is always ignored
    - o. Direct toward recursion
    - p. Should be maximized
- Indicates the duration of an activity page 106
  - 427. The prefix code generated by Huffman algorithm -----the expected length of the encoded string.
- Minimizes page 102
- Balances
- Maximizes
- Keeps constant
  - 428. An optimization problem is one in which you want to find,
    - ► Not a solution
    - ► An algorithm
    - ► Good solution
    - ► The best solution (Page 97)
  - 429. Although it requires more complicated data structures, Prim's algorithm for a minimum spanning tree is better than Kruskal's when the graph has a large number of vertices.
    - **►** True
    - ► False
  - 430. If a problem is in NP, it must also be in P.
    - ► True
    - ► False

- ► unknown (Page 173)
- 431. What is generally true of Adjacency List and Adjacency Matrix representations of graphs?
  - ► Lists require less space than matrices but take longer to find the weight of an edge (v1,v2)
  - ightharpoonup Lists require less space than matrices and they are faster to find the weight of an edge (v1,v2)
  - ► Lists require more space than matrices and they take longer to find the weight of an edge (v1,v2)
  - Lists require more space than matrices but are faster to find the weight of an edge (v1,v2)
- 432. If a graph has v vertices and e edges then to obtain a spanning tree we have to delete
  - ➤ v edges.
  - $\triangleright$  v e + 5 edges
  - $\triangleright$  v + e edges.
  - ► None of these
- 433. Maximum number of vertices in a Directed Graph may be |V2 |
  - ► True
  - **►** False
- 434. The Huffman algorithm finds a (n) \_\_\_\_\_solution.
  - **▶** Optimal
  - ► Non-optimal
  - ► Exponential
  - ► Polynomial
- 435. The Huffman algorithm finds an exponential solution
  - ► True
  - ► False
- 436. The greedy part of the Huffman encoding algorithm is to first find two nodes with larger frequency.
  - ► True
  - ► False (Page 100)
- 437. one The codeword assigned to characters by the Huffman algorithm have the property that no codeword is the postfix of any other.
  - ► True (Page 101)
  - ► False
- 438. Huffman algorithm uses a greedy approach to generate a postfix code T that minimizes the expected length B (T) of the encoded string.
  - ► True
  - ► False (Page 102)

- 439. Shortest path problems can be solved efficiently by modeling the road map as a graph.
  - ► True (Page 153)
  - ► False
- 440. Dijkestra's single source shortest path algorithm works if all edges weights are non-negative and there are negative cost cycles.
  - ► True
  - ► False (Page 159)
- 441. Bellman-Ford allows negative weights edges and negative cost cycles.
  - ► True
  - ► False (Page 159)
- 442. The term "coloring" came form the original application which was in architectural design.
  - ► True
  - ► False (Page 176)
- 443. In the clique cover problem, for two vertices to be in the same group, they must be adjacent to each other.
  - **►** True (Page 176)
  - ► False
- 444. Dijkstra's algorithm is operates by maintaining a subset of vertices
  - ► True (Page 155)
  - ► False
- 445. The difference between Prim's algorithm and Dijkstra's algorithm is that Dijkstra's algorithm uses a different key.
  - ► True ( Page 156)
  - ► False
- 446. We do sorting to,
  - ► keep elements in random positions
  - ► keep the algorithm run in linear order
  - ▶ keep the algorithm run in (log n) order
  - keep elements in increasing or decreasing order (Page 40)
- 447. After partitioning array in Quick sort, pivot is placed in a position such that
  - ► Values smaller than pivot are on left and larger than pivot are on right
  - (Page 48) ► Values larger than pivot are on left and smaller than pivot are on right
  - ▶ Pivot is the first element of array
  - ▶ Pivot is the last element of array
- 448. Merge sort is stable sort, but not an in-place algorithm
  - True (Page 54)
  - ► False

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449.	In counting sort, once we know the ranks, we simplynumbers to
	their final positions in an output array.
	▶ Delete
	copy (Page 57)
	► Mark
	► arrange
450.	one Dynamic programming algorithms need to store the results of intermediate
	sub-problems.
	► True (Page 75)
-05	► False
451.	A p $\times$ q matrix A can be multiplied with a q $\times$ r matrix B. The result will be a p
1)	× r matrix C. There are (p . r) total entries in C and each takesto
U.	compute.
	► O (q) (Page 84)
	► O(1)
	► O (n 2)
450	► O (n3) -
452.	is a graphical representation of an algorithm
	$\Sigma$ notation
	Θ notation
	Flowchart
450	Asymptotic notation
455.	one Which of the following is calculated with big o notation?
	• Lower bounds
	• Upper bounds (Page 25)
	• ▶ Both upper and lower bound
454	• Medium bounds
454.	Merge sort makes two recursive calls. Which statement is true after these
	recursive calls finish, but before the merge step?
-	• The array elements form a heap
	• Elements in each half of the array are sorted amongst themselves
	• Elements in the first half of the array are less than or equal to elements in
	the second half of the array
455	None of the above 7
455.	Who invented Quick sort procedure?
	• Hoare
	• Sedgewick
	• Mellroy
	• ► Coreman

456. What is the solution to the recurrence T(n) = T(n/2) + n, T(1) = 1

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		• ►O(logn)
		• ►O(n) (Page 37)
		• ►O(nlogn)
		• ► O(2n)
	457.	one Consider the following Huffman Tree The binary code for the string TEA
	10 7 1	is
		• ► 10 00 010
		• ►011 00 010
		• ► 10 00 110
		• ►11 10 110
	158	A greedy algorithm does not work in phases.
1	150.	True
11	U.	► False (Page 97)
	459.	Can an adjacency matrix for a directed graph ever not be square in shape?
		• ►Yes
Y		• ▶No
Y /	460.	One of the clever aspects of heaps is that they can be stored in arrays without
		using any
		● Pointers (Page 40)
		constants
		• <b>&gt;</b> variables
		• Functions
	461.	Merge sort requires extra array storage, •
		► True (Page 54) •
		► False
	462.	Non-optimal or greedy algorithm for money change takes
		• ► O(k) (Page 99)
	- 5	• ►O(kN)
		• ► O(2k)
		• ►O(N)
	463.	The Huffman codes provide a method of encoding data inefficiently when
		coded using ASCII standard.
		• ► True
		• ► False (Page 99)
	464.	Using ASCII standard the string abacdaacac will be encoded with
		hits

• ►80 (Page 99)

▶160▶320

- ▶100
- 465. Using ASCII standard the string abacdaacac will be encoded with 160 bits.
  - ► True
  - False (Page 99)
- 466. one Using ASCII standard the string abacdaacac will be encoded with 320 bits.
  - ► True
  - False (Page 99)
- 467. Using ASCII standard the string abacdaacac will be encoded with 100 bits.
  - ► True
  - False (Page 99)
- 468. Using ASCII standard the string abacdaacac will be encoded with 32 bytes
  - ► True
  - ► False (Page 99)
- 469. The greedy part of the Huffman encoding algorithm is to first find two nodes with smallest frequency.
  - ► True (Page 100)
  - ►False
- 470. The greedy part of the Huffman encoding algorithm is to first find two nodes with character frequency
  - ►True
  - ► False (Page 100)
- 471. Huffman algorithm uses a greedy approach to generate an antefix code T that minimizes the expected length B (T) of the encoded string.
  - ►True
  - False (Page 102)
- 472. Depth first search is shortest path algorithm that works on un-weighted graphs.
  - ► True
  - ► False (Page 153)
- 473. one Dijkestra s single source shortest path algorithm works if all edges weights are non negative and there are no negative cost cycles.
  - **True** (Page 159)
  - ►False
- 474. Dijkestra s single source shortest path algorithm works if all edges weights are negative and there are no negative cost cycles.
  - ► True •
  - ► False (Page 159
- 475. Floyd-Warshall algorithm is a dynamic programming algorithm; the genius of the algorithm is in the clever recursive formulation of the shortest path problem.

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476.	• ► True (Page 162) • ► Flase Floyd-Warshall algorithm, as in the case with DP algorithms, we avoid
	recursive evaluation by generating a table for
	• <u>▶</u> k
	• ► k ij d (Page 164)
	• ►True
	• ► Flase
477.	The term coloring came from the original application which was in map
	drawing.
10	• ► True (Page 176)
11	• ► False
478.	In the clique cover problem, for two vertices to be in the same group, they must
$\vee_{A}$	beeach other.
	• Apart from
	• ► Far from
	• ► Near to
4=0	• Adjacent to (Page 176)
479.	Fixed-length codes may not be efficient from the perspective ofthe
	total quantity of data. Select correct option:
	► Minimizing (Page 99)  Averaging
	➤ Averaging  ➤ Maximizing
	► Summing
480	In greedy algorithm, at each phase, you take theyou can get right
100.	now, without regard for future consequences.
	► Worst
	► Minimum
	► Good
- 1	► Best (Page 97)
481.	one The difference between Prim s algorithm and Dijkstra s algorithm is that
	Dijkstra s algorithm uses a same key.
	• ► True
	• ► False ( Page 156)
482.	If a problem is in NP-complete, it must also be in NP.
	► True (Page 178)
	► False
483.	If there are n items, there arepossible combinations of the items.

► 2^n (Page 92)

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	▶ 3^n
484.	Using ASCII code, each character is represented by a fixed-length code word
	ofbits per character.
	▶4
	<b>▶</b> 6
	▶ <mark>8 (P)age 99)</mark>
	<b>▶</b> 10
485.	In Knapsack Problem, the thief's goal is to put items in the bag such that the
	of the items does not exceed the limit of the bag.
-	► <mark>Value (Page 91)</mark>
16	► Weight
11	► Length
	► Balance
486.	one The knapsack problem does not belong to the domain of optimization
	problems. ▶ True
	► False (Page 91)
487.	In Huffman encoding, for a given message string, the frequency of occurrence
	(relative probability) of each character in the message is determined last.
	True
	► False (Page 100)
488.	Fixed-length codes are known for easy break up of a string into its individual
	characters.
	► True (Page 99)
	▶ False
489.	Question No: 8 (Marks: 1) - Please choose one InKnapsack Problem,
	limitation is that an item can either be put in the bag or not-fractional items are
	not allowed.
1	1 1 1 2 1 3 1 3 1 4 - 1 6 5 9 2 9 4 4
	▶ <mark>0/1 (Page 91)</mark>
400	Fractional
490.	one The term "coloring" came from the original application which was in
	architectural design.
	True
401	False (Page173)
491.	In Knapsack Problem, value and weight both are to be under consideration.
	True (page 91)
400	False
492.	Time complexity of DP based algorithm for computing the minimum cost of
	chain matrix Multiplication is .

	▶ log n
	▶n
	<b>▶</b> n2
	▶ n3 (Page 90)
493.	Qestion No: 12 (Marks: 1) - Please choose one In DP based solution of
	knapsack problem, to compute entries of V we will imply a/an
	approach.
	➤ Subjective
	► Inductive (Page 93)
- 2	▶ Brute force
10	► Combination
494.	A greedy algorithm sometimes works well for optimization problems.
// ) :	► True (Page 97)
	► False
495.	In Huffman encoding, frequency of each character can be determined by
/	parsing the message andhow many times each character (or
	symbol) appears. Printing
	► Incrementing
	► Counting (Page 100)
	▶ Deleting
496.	Greedy algorithm can do very poorly for some problems.
	► True (Page 97)
	► False 14
497.	one The Huffman codes provide a method of data efficiently.
	► Reading
	► Encoding (Page 99)
	▶ Decoding
	▶ Printing
498.	Inbased solution of knapsack problem, we consider 2 cases, Leave
	object Or Take object.
	object Or Take object.  ► Brute force  ► Dynamic programming (Page 93)
	► Dynamic programming (Page 93)
499.	Those problems in which Greedy finds good, but not always best is called a
	greedy
	► Algorithm
	► Solution
	► Heuristic (Page 97)
	► Result
500.	one In brute force based solution of knapsack problem, we consider 2 cases,
	Leave object Or Take object.

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		▶TRUE
		► FALSE (Page 97)
	501.	
		► Minimization
		► Averaging
		► Optimization (Page 97)
		► Maximization
	502.	Using ASCII standard the string abacdaacac will be encoded with 10 bytes.
		• ► True (Page 101)
	- 1	• ► False
	503.	Inalgorithm, you hope that by choosing a local optimum at each step,
1	11	you will end up at a global optimum.
11		► Simple
		▶ Non Greedy
( )		► Greedy (Page 97)
V		▶ Brute force
y' /	504.	one Huffman algorithm uses a greedy approach to generate an prefix code T
- /		that minimizes the expected length B (T) of the encoded string.
		• ► True (Page 102)
		• ► False
	505.	Counting Money problem is an example which cannot be optimally solved by
		greedy algorithm.
		► True (Page 97)
		► False
	506.	Huffman algorithm generates an optimum prefix code.
		► True (Page 102)
		► False
	507.	If the string "lmncde" is coded with ASCII code, the message length would be
	-	bits.
	- 4	<b>▶</b> 24
		<ul> <li>≥ 24</li> <li>≥ 36</li> <li>≥ 48 (6*8=48)</li> <li>≥ 60</li> </ul>
		► 48 (6*8=48)
		▶60 VIII mghel
	508.	
		minimum cost of chain matrix multiplication.
		►2
		► 3 (Page 90)
		►4 ►5
		<b>▶</b> 5

- 509. Inductive approach to compute entries of V is implied in \_\_\_\_\_\_based solution of knapsack problem.

  ▶ Brute force

  ▶ Dynamic programming (Page 93)
- 510. A number of lectures are to be given in a single lecture hall. Optimum scheduling for this is an example of Activity selection.
  - ► True (Page 105)
  - ► False
- 511. The activity scheduling is a simple scheduling problem for which the greedy algorithm approach provides a/an solution.
  - **►** Simple
  - ► Sub optimal
  - ► Optimal (Page 105)
  - ► Non optimal
- 512. Question # 1 of 10 (Marks: 1) Please choose one The string |xyz|, if coded with ASCII code, the message length would be 24 bits.
  - **►** True
  - ► False
- 513. An application problem is one in which you want to find, not just a solution, but the solution.
  - ► Simple
  - ► Good (Page 113)
  - ▶ Best
  - ► Worst
- 514. A dense undirected graph is:
  - ► A graph in which  $E = O(V^2)$  click here 4 detail
  - ightharpoonup A graph in which E = O(V)
  - ightharpoonup A graph in which  $E = O(\log V)$
  - ► All items above may be used to characterize a dense undirected graph
- 515. Suppose that a graph G = (V,E) is implemented using adjacency lists. What is the complexity of a breadth-first traversal of G?
  - ► O(|V |^2)
  - ► O(|V | |E|)
  - ► O(|V |^2|E|)
  - ightharpoonup O(|V| + |E|) pg 116
- 516. Which is true statement?
  - ► Breadth first search is shortest path algorithm that works on un-weighted graphs (Page 153)
  - ▶ Depth first search is shortest path algorithm that works on un-weighted graphs. ▶ Both of above are true.

- ▶ None of above are true.
- 517. one Forward edge is:
  - ► (u, v) where u is a proper descendent of v in the tree.
  - (u, v) where v is a proper descendent of u in the tree. (Page 129)
  - ► (u, v) where v is a proper ancesstor of u in the tree.
  - ► (u, v) where u is a proper ancesstor of v in the tree.
- 518. What general property of the list indicates that the graph has an isolated vertex? ► There is Null pointer at the end of list.
  - ► The Isolated vertex is not handled in list.
  - ▶ Only one value is entered in the list.
  - ► There is at least one null list.
- 519. one If you find yourself in maze the better traversal approach will be:
  - **▶** BFS
  - ▶ BFS and DFS both are valid
  - ► Level order
  - **▶** DFS
- 520. In digraph G=(V,E); G has cycle if and only if
  - ► The DFS forest has forward edge.
  - ► The DFS forest has back edge (Page 131)
  - ► The DFS forest has both back and forward edge
  - ► BFS forest has forward edge
- 521. Back edge is:
  - (u, v) where v is an ancestor of u in the tree. (Page 128)
  - $\blacktriangleright$  (u,v) where u is an ancesstor of v in the tree.
  - ► (u, v) where v is an predcessor of u in the tree.
  - ▶ None of above
- 522. Which statement is true?
  - ► If a dynamic-programming problem satisfies the optimal-substructure property, then a locally optimal solution is globally optimal.
  - ► If a greedy choice property satisfies the optimal-substructure property, then a locally optimal solution is globally optimal.
  - ▶ Both of above
  - ► None of above
- 523. Cross edge is:
  - ► (u, v) where u and v are not ancestor of one another
  - (u, v) where u is ancesstor of v and v is not descendent of u.
  - (u, v) where u and v are not ancestor or descendent of one another (Page 129)
  - (u, v) where u and v are either ancestor or descendent of one another.

- 524. Kruskal's algorithm (choose best non-cycle edge) is better than Prim's (choose best tree edge) when the graph has relatively few edges.
  - ► True click here 4 detail
  - ► False
- 525. Which is true statement in the following?
  - ► Kruskal algorithm is multiple source technique for finding MST. click here for detail ► Kruskal's algorithm is used to find minimum spanning tree of a graph, time complexity of this algorithm is O(EV)
  - ▶ Both of above
  - ► Kruskal's algorithm (choose best non-cycle edge) is better than Prim's (choose best Tree edge) when the graph has relatively few edges. click here 4 detail
- 526. Kruskal's algorithm (choose best non-cycle edge) is better than Prim's (choose best tree edge) when the graph has relatively few edges.
  - ► True click here 4 detail
  - ► False
- 527. What algorithm technique is used in the implementation of Kruskal solution for the MST?
  - ► Greedy Technique (Page 142)
  - ► Divide-and-Conquer Technique
  - ► Dynamic Programming Technique
  - ► The algorithm combines more than one of the above techniques
- 528. What is the time complexity to extract a vertex from the priority queue in Prim's algorithm?
  - **▶** O (log E)
  - **►** (V)
  - ► (V+E)
  - ► O (log V) (Page 152)
- 529. The relationship between number of back edges and number of cycles in DFS is, ▶ Both are equal
  - ► Back edges are half of cycles
  - ► Back edges are one quarter of cycles
  - There is no relationship between no. of edges and cycles (Page 131)
- 530. You have an adjacency list for G, what is the time complexity to compute Graph transpose G^T.?
  - $\triangleright$  (V + E) (Page 138)
  - **►** (V E)
  - **▶** (V)
- 531. A digraph is strongly connected under what condition?

- ► A digraph is strongly connected if for every pair of vertices u, v e V, u can reach v. ► A digraph is strongly connected if for every pair of vertices u, v e V, u can reach v and vice versa. (Page 135)
- ► A digraph is strongly connected if for at least one pair of vertex u, v e V, u can reach v and vice versa.
- ► A digraph is strongly connected if at least one third pair of vertices u, v e V, u can reach v and vice versa.
- 532. In in-place sorting algorithm is one that uses arrays for storage:
  - ► An additional array
  - ► No additional array (Page 54)
  - ▶ Both of above may be true according to algorithm
  - ► More than 3 arrays of one dimension.
- 533. In stable sorting algorithm
  - ▶ One array is uses
  - ► In which duplicating elements are not handled.
  - ► More then one arrays are required.
  - ► Duplicating elements remain in same relative position after sorting. (Page 54)
- 534. Which sorting algorithm is faster:
  - $ightharpoonup O(n^2)$
  - ► O(nlogn) (Page 46)
  - ightharpoonup O(n+k)
  - $ightharpoonup O(n^3)$
- 535. In Quick sort algorithm, constants hidden in T(n lg n) are
  - ► Large
  - ► Medium
  - ▶ Not known
  - **►** Small
- 536. Quick sort is based on divide and conquer paradigm; we divide the problem on base of pivot element and:
  - ▶ There is explicit combine process as well to conquer the solution.
  - ▶ No work is needed to combine the sub-arrays, the array is already sorted
  - ► Merging the sub arrays
  - ► None of above. (Page 51
- 537. Dijkstra's algorithm:
  - ► Has greedy approach to find all shortest paths
  - ► Has both greedy and Dynamic approach to find all shortest paths
  - ► Has greedy approach to compute single source shortest paths to all other vertices (Page 154)

- ► Has both greedy and dynamic approach to compute single source shortest paths to all other vertices.
- 538. Which may be stable sort:
  - ▶ Bubble sort
  - ► Insertion sort
  - ► Both of above (page 54)
  - ► Selection sort
- 539. Question # 1 of 10 (Marks: 1) Please choose one In the analysis of Selection algorithm, we eliminate a constant fraction of the array with each phase; we get the convergent series in the analysis,
  - ► linear
  - **▶** arithmetic
  - ▶ geometric (page 37)
  - **▶** exponent
- 540. How much time merge sort takes for an array of numbers?
  - ► T(n^2)
  - ► T(n) (Page 40)
  - ► T( log n)
  - ightharpoonup T(n log n)
- 541. Counting sort has time complexity:
  - $\triangleright$  O(n)  $\triangleright$  O(n+k)
  - **▶** O(k)
  - ► O(nlogn)
- 542. The analysis of Selection algorithm shows the total running time is indeed in no
  - **▶** arithmetic
  - **▶** geometric
  - ► linear (Page 37)
  - **▶** orthogonal
- 543. Sorting is one of the few problems where provable \_\_\_\_\_bonds exits on how fast we can sort,
  - **▶** upper
  - ► lower (Page 39)
  - ▶ average
  - ▶ log n
- 544. In the analysis of Selection algorithm, we make a number of passes, in fact it could be as many as
  - **►** T(n)
  - ightharpoonup T(n/2)
  - $\triangleright$  log n (Page 37)

- $\triangleright$  n / 2 + n / 4
- 545. number of nodes in a complete binary tree of height h is
  - $ightharpoonup 2^{(h+1)} 1 \text{ (Page 40)}$
  - $\triangleright 2 * (h+1) 1$
  - ► 2 \* (h+1)
  - $\blacktriangleright$  ((h+1)  $^{\land}$  2) 1 25
- 546. How many elements do we eliminate in each time for the Analysis of Selection algorithm?
  - ► n / 2 elements (Page 37)
  - $\blacktriangleright$  (n / 2) + n elements
  - ▶ n / 4 elements
  - ▶ 2 n elements
- 547. Slow sorting algorithms run in,
  - ightharpoonup T(n^2) (Page 39)
  - **►** T(n)
  - **►** T( log n)
  - ightharpoonup T(n log n)
- 548. Counting sort is suitable to sort the elements in range 1 to k:
  - ► K is large
  - ► K is small (Page 57)
  - ► K may be large or small
  - ► None
- 549. Heaps can be stored in arrays without using any pointers; this is due to the \_\_\_\_\_\_nature of the binary tree,
  - ► left-complete (Page 40)
  - ► right-complete
  - ► tree nodes
  - ► tree leaves
- 550. Sieve Technique can be applied to selection problem?
  - ► True (Page 35)
  - ► False
- 551. A heap is a left-complete binary tree that conforms to the
  - ▶ increasing order only
  - ▶ decreasing order only
  - ► heap order (Page 40)
  - ► (log n) order
- 552. Divide-and-conquer as breaking the problem into a small number of
  - **▶** pivot
  - **▶** Sieve
  - ► smaller sub problems (Page 34)

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		► Selection
	553.	In Sieve Technique we do not know which item is of interest
		True (Page 34)
		► False
	554.	The recurrence relation of Tower of Hanoi is given below $T(n)=\{1 \text{ if } n=1 \text{ and } 2\pi(-1): (n-1) \text{ if } n=1 \text{ and }$
		2T(n-1) if n >1 In order to move a tower of 5 rings from one peg to another,
		how many ring moves are required?
		► 16
		► 10
	-05	<b>▶</b> 32
- 2	17	► 31 Click here 4 detail
×1	555.	one For the heap sort, access to nodes involves simple
11	U,	operations.
/		▶ arithmetic (Page 41)
		▶ binary
V A		▶ algebraic
Y /		▶ logarithmic
	556.	For the sieve technique we solve the problem,
		recursively (Page 34)
		► mathematically
		▶ precisely
		► accurately
	557.	The sieve technique works in as follows
		▶ phases (Page 34)
		▶ numbers
		▶ integers
	N.	▶ routines
	558.	A (an)is a left-complete binary tree that conforms to the heap order
	1	► heap (Page 40)
		▶ binary tree
		<ul><li>▶ binary search tree</li><li>▶ array</li></ul>
	559.	
		>5 CHITISTIC ST
		▶ many
		▶ 1 (Page 34)

560. Analysis of Selection algorithm ends up with,

**►** T(n)

► T(1/1+n)

- ightharpoonup T(n/2)
- ightharpoonup T((n/2) + n) (Page 37)
- 561. For the heap sort we store the tree nodes in
  - ► level-order traversal (Page 40)
  - ▶ in-order traversal
  - ▶ pre-order traversal
  - ▶ post-order traversal
- 562. The reason for introducing Sieve Technique algorithm is that it illustrates a very important special case of,
  - ▶ divide-and-conquer (Page 34)
  - ► decrease and conquer
  - ► greedy nature
  - ► 2-dimension Maxima
- 563. Theta asymptotic notation for T (n):
  - ▶ Set of functions described by: c1g(n)Set of functions described by c1g(n)>=f(n) for c1 s
  - ► Theta for T(n)is actually upper and worst case comp
  - ► Set of functions described by:
  - ightharpoonup clg(n)
- 564. Sieve Technique applies to problems where we are interested in finding a single item from a larger set of \_\_\_\_\_\_
  - ▶ n items (Page 34)
  - **▶** phases
  - **▶** pointers
  - **▶** constant
- 565. Memorization is?
  - ► To store previous results for future use
  - ► To avoid this unnecessary repetitions by writing down the results of recursive calls and looking them up again if we need them later (Page 47)
  - ► To make the process accurate
  - ▶ None of the above
- 566. Quick sort is
  - ► Stable & in place
  - Not stable but in place (Page 57)
  - ► Stable but not in place
  - ► Some time stable & some times in place
- 567. One example of in place but not stable algorithm is
  - ► Merger Sort
  - ► Quick Sort (Page 54)
  - ► Continuation Sort

- ▶ Bubble Sort
- 568. Continuation sort is suitable to sort the elements in range 1 to k
  - ► K is Large
  - ► K is not known
  - ► K may be small or large
  - ► K is small (Page 57)
- 569. Which may be a stable sort?
  - ► Merger
  - ► Insertion
  - ► Both above (Page 54)
  - ▶ None of the above
- 570. An in place sorting algorithm is one that uses arrays for storage
  - ► Two dimensional arrays
  - ► More than one array
  - ► No Additional Array (Page 54)
  - ▶ None of the above
- 571. Continuing sort has time complexity of?
  - ightharpoonup O(n)
  - ightharpoonup O(n+k)
  - ► O(nlogn)
  - **▶** O(k)
- 572. single item from a larger set of \_\_\_\_
  - ▶ n items (Page 34)
  - **▶** phases
  - **▶** pointers
  - **▶** vconstant
- 573. For the Sieve Technique we take time
  - ► T(nk) ( Page 34)
  - ightharpoonup T(n/3)
  - ► n^2
  - ► n/3
- 574. One Example of in place but not stable sort is
  - ► Quick (Page 54)
  - ► Heap
  - ► Merge
  - **▶** Bubble
- 575. Consider the following Algorithm: Factorial (n) { if (n=1) return 1 else return (n \* Factorial(n-1)) } Recurrence for the following algorithm is:
  - ightharpoonup T(n) = T(n-1) + 1
  - ightharpoonup T(n) = nT(n-1) + 1

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		T(n) = T(n-1) + n
		ightharpoonup T(n) = T(n(n-1)) + 1
	576.	Due to left complete nature of binary tree, the heap can be stored in
		► Arrays (Page 40)
		► Structures
		► Link Lis
		► Stack
	577.	What type of instructions Random Access Machine (RAM) can execute?
		► Algebraic and logic
	-	► Geometric and arithmetic
	16	► Arithmetic and logic (Page 10)
1	11	► Parallel and recursive
11	578.	What is the total time to heapify?
	- 4	O(log n) (Page 43)
( )		ightharpoonup O(n log n)
V		$ ightharpoonup O(n2 \log n)$
y' /		► O(log2 n)
· /	579.	word Algorithm comes from the name of the muslim author
		► Abu Ja'far Mohammad ibn Musa al-Khowarizmi.
	580.	al-Khwarizmi's work was written in a book titled

- ▶ al Kitab al-mukhatasar fi hisab al-jabr wa'l-muqabalah
- 581. Random access machine or RAM is a/an
  - ► Machine build by Al-Khwarizmi
  - ► Mechanical machine
  - ► Electronics machine
  - ► Mathematical model (Page 10)
- 582. A RAM is an idealized machine with\_ random-access memory.
  - ► 256MB
  - ► 512MB
  - ► an infinitely large (Page 10)
  - ► 100GB
- 583. What will be the total number of max comparisons if we run brute-force maxima algorithm with n elements?
  - ▶ 2 n
  - ► 2 n n
  - ► n (Page 14)
- 584. Due to left complete nature of binary tree, the heap can be stored in
  - ► Arrays (Page 40)

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	<ul><li>▶ Structures</li><li>▶ Link Lis</li><li>▶ Stack</li></ul>
585.	one What type of instructions Random Access Machine (RAM) can execute?
	<ul><li>▶ Algebraic and logic</li><li>▶ Geometric and arithmetic</li></ul>
	► Arithmetic and logic (Page 10)  ► Parallel and recursive
586.	What is the total time to be wift?
17	<ul> <li>Vital is the total time to heapily?</li> <li>► O(log n) (Page 43)</li> <li>► O(n log n)</li> <li>► O(n2 log n)</li> </ul>
))	<ul><li>▶ O(n2 log n)</li><li>▶ O(log2 n)</li></ul>
587.	Random access machine or RAM is a/an
	► Machine build by Al-Khwarizmi
	► Mechanical machine
	► Electronics machine
	► Mathematical model (Page 10)
588.	A RAM is an idealized machine withrandom-access
	memory.
	► 256MB
	► 512MB
	► an infinitely large (Page 10)
	▶ 100GB
589.	What will be the total number of max comparisons if we run brute-force
	maxima algorithm with n elements?
	► 2 n
	≥ 2 n n
- 1	$ \underset{8 \text{ n}}{\triangleright} \underset{8 \text{ n}}{(Page 14)} 0304 - 1659294 $
590	The sieve technique works where we have to finditem(s) from a
570.	
	large input.  ► Single (Page 34)  ► Two
	► Two
	► Three
	► Similar
591.	In which order we can sort?
	► increasing order only
	► decreasing order only
	585. 586. 587.

► increasing order or decreasing order (Page 39)

	1 .1	•		. •
$\sim$	hoth	at the	came	time
	17071.11	at the	Same	

592.	For t	he	hean	sort	we	store	the	tree	nodes	in
5,2.	1 01 6	110	ncap	BOIL	** C	Stor C	uic	ucc	Houcs	111

- ► level-order traversal (Page 40)
- ▶ in-order traversal
- ▶ pre-order traversal
- ▶ post-order traversal
- 593. In the analysis of Selection algorithm, we eliminate a constant fraction of the array with each phase; we get the convergent \_\_\_\_\_\_\_ series in the analysis, ▶linear
  - **▶** arithmetic
  - ▶ geometric (Page 37)
  - **▶** exponent
- 594. How much time merge sort takes for an array of numbers?
  - ► T(n^2)
  - **►** T(n)
  - ► T( log n)
  - ightharpoonup T(n log n) (Page 40)
- 595. Memoization is?
  - ► To store previous results for future use
  - ► To avoid this unnecessary repetitions by writing down the results of recursive calls and looking them up again if we need them later (page 74)
  - ► To make the process accurate
  - ▶ None of the above
- 596. Cont sort is suitable to sort the elements in range 1 to k
  - ► K is Large
  - ► K is not known
  - ► K may be small or large
  - ► K is small (Page 57)
- 597. In place stable sorting algorithm.
  - ▶ If duplicate elements remain in the same relative position after sorting (Page
  - 54) ► One array is used
  - ► More than one arrays are required
  - ► Duplicating elements not handled
- 598. Sorting is one of the few problems where provable bonds exits on how fast we can sort, ▶upper ▶lower (Page 39) ▶average ▶log n
- 599. Question No: 65 (Marks: 1) Please choose one Counting sort has time complexity: ►O(n) (Page 58)
  - ightharpoonup O(n+k)
  - **►** O(k)
  - ► O(nlogn)

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600	). The running time of quick sort depends heavily on the selection of
	► No of inputs
	► Arrangement of elements in array
	► Size o elements
	► Pivot elements (Page 49)
601	. Which may be stable sort:
	▶ Bubble sort
	► Insertion sort
	► Both of above (Page 54)
602	2. In Quick Sort Constants hidden in T(n log n) are
17	► Large
1	► Medium
	► Small
1	▶ Not Known
603	6. Quick sort is based on divide and conquer paradigm; we divide the problem
	base of pivot element and:
	► There is explicit combine process as well to conquer the solution.
	▶ No work is needed to combine the sub-arrays, the array is already sorted
	► Merging the sub arrays
60.	None of above. (Page 51)
604	A point p in 2-dimensional space is usually given by its integer
	coordinate(s)
	p.x only
	p.y only
	• p.x & p.z
605	p.x & p.y (Page 10)
003	6. In we have to find rank of an element from given input.
	► Merge sort algorithm  ► Selection problem (Page 34)
	► Brute force technique
606	<ul> <li>▶ Plane Sweep algorithm</li> <li>In Heap Sort algorithm, if heap property is violated</li> <li>▶ We call Build heap procedure</li> </ul>
000	➤ We call Build heap procedure
	► We call Heapify procedure  ■ We call Heapify procedure
	► We ignore
	► Heap property can never be violated
607	7. Upper bound requires that there exist positive constants c2 and n0 such that
237	$f(n)$ c2n for all $n \le n0$ (ye question ghalat lag raha hai mujhae
	► Less than

► Equal to or Less than (Page 25)

- ► Equal or Greater than
- ► Greater than
- 608. A RAM is an idealized algorithm with takes an infinitely large random-access memory. ► True
  - ► False (Page 10)
- 609. \_is one of the few problems, where provable lower bounds exist onhow fast we can sort.
  - **▶** Searching
  - **►** Sorting
  - ▶ Both Searching & Sorting
  - **▶** Graphing
- 610. Floor and ceiling are to calculate while analyzing algorithms.
  - ► Very easy
  - ► Usually considered difficult (Page 31)
- 611. In Heap Sort algorithm, the maximum levels an element can move upward is
  - ► Theta (log n) (Page 43)
  - ➤ Order (log n)
  - ► Omega (log n)
  - ▶ O (1) i.e. Constant time
- 612. A point p in 2-dimensional space is usually given by its integer coordinate(s)
  - ▶ p.x only p.y
  - only p.x & p.z
  - ▶ p.x & p.y (Page 17)
- 613. In Heap Sort algorithm, the total running time for Heapify procedure is
  - ► Theta (log n) (Page 43)
  - ► Order (log n)
  - ► Omega (log n)
  - ► O (1) i.e. Constant time
- 614. Algorithm is a mathematical entity, which is independent of a specific machine and operating system.
  - ► True
  - ► False (Page 7)
- 615. While Sorting, the ordered domain means for any two input elements x and y satisfies only.
  - $\triangleright$  x < y
  - $\triangleright x > y$
  - $\triangleright x = y$

- ► All of the above (Page 39)
- 616. Quick sort is best from the perspective of Locality of reference.
  - ► True (Page 9)
  - ► False
- 617. In Heap Sort algorithm, we build for ascending sort.
  - ► Max heap (Page 41)
  - ► Min heap
- 618. In Sieve Technique, we know the item of interest.
  - ► True
  - ► False (Page 34)
- 619. While solving Selection problem, in Sieve technique we partition input data
  - ► In increasing order
  - ► In decreasing order
  - ► According to Pivot (Page 35)
  - **►** Randomly
- 620. In pseudo code, the level of details depends on intended audience of the algorithm.
  - ► True (Page 12)
  - ► False
- 621. If the indices passed to merge sort algorithm are\_\_\_\_\_, then this means that there is only one element to sort. ▶ Small
  - ► Large
  - ► Equal (Page 28)
  - ► Not Equal

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